

Data Structures – CST 201

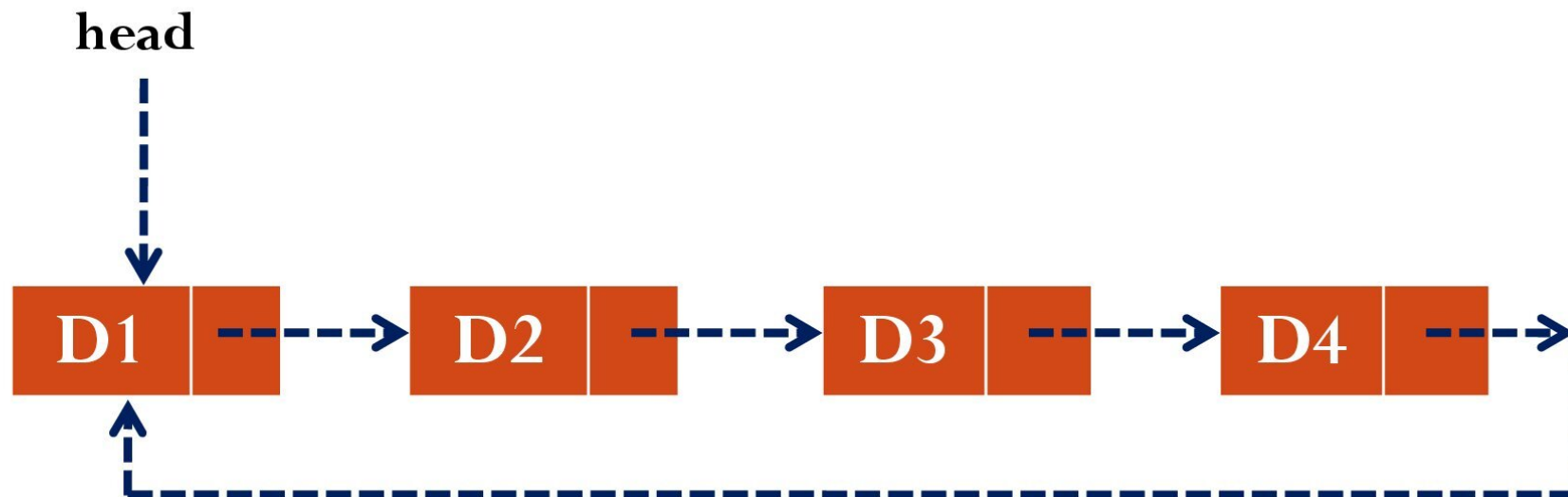
Module ~ 3

Syllabus

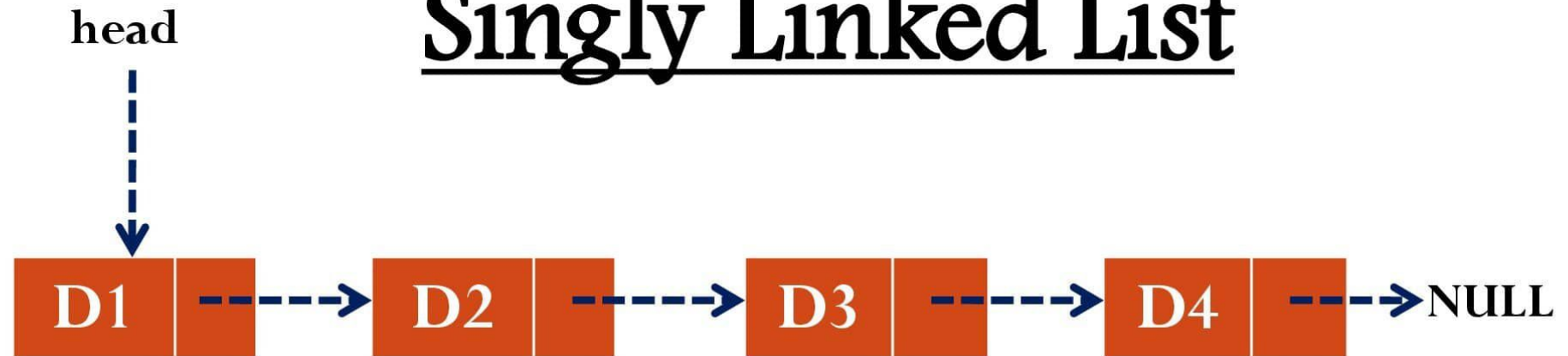
- **Linked List and Memory Management**
 - Self Referential Structures
 - Dynamic Memory Allocation
 - Singly Linked List~Operations on Linked List.
 - Doubly Linked List
 - **Circular Linked List**
 - Stacks using Linked List
 - Queues using Linked List
 - Polynomial representation using Linked List
 - Memory allocation and de~allocation
 - First~fit, Best~fit and Worst~fit allocation schemes

Circular Linked List

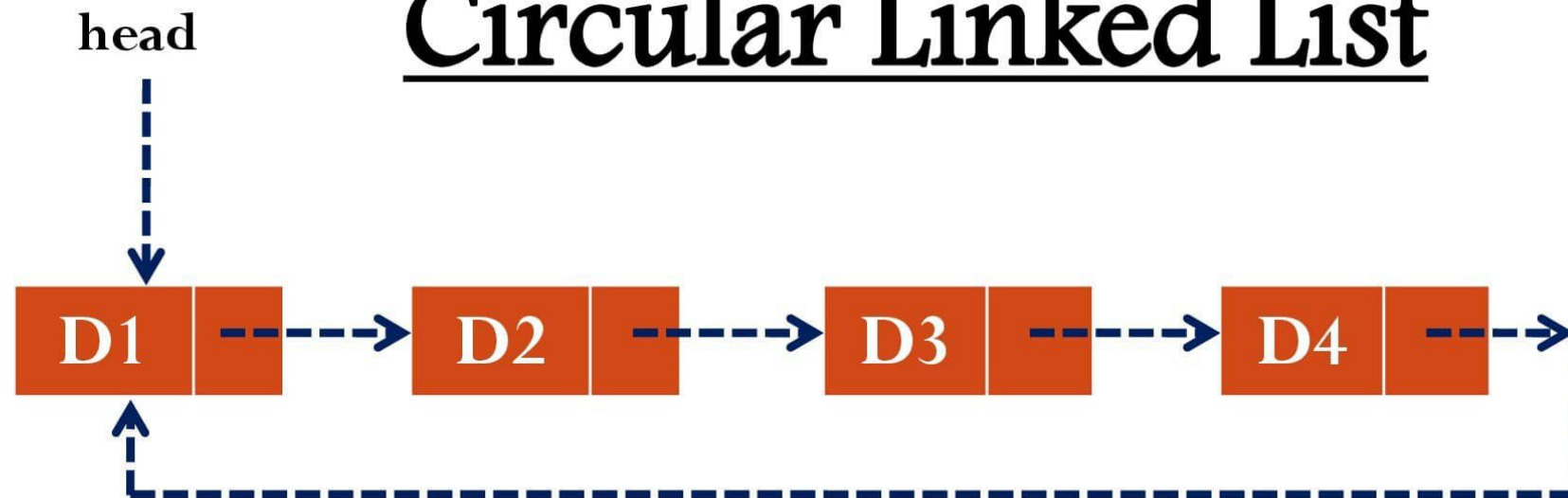
- In a single linked list, the link field of the last node is null.
- A linked list with last node points to the first node is called **circular linked list**



Singly Linked List



Circular Linked List



Circular Linked List

- **Advantages**

- **Accessibility of a member node in the list:** In an ordinary list, a member node is accessible from a particular node, that is, from the header node only. But in the circular linked list, every member node is accessible from any node
- **Avoided Null link problem:** Null value in link field may create problem during the execution of the program if proper care is not taken

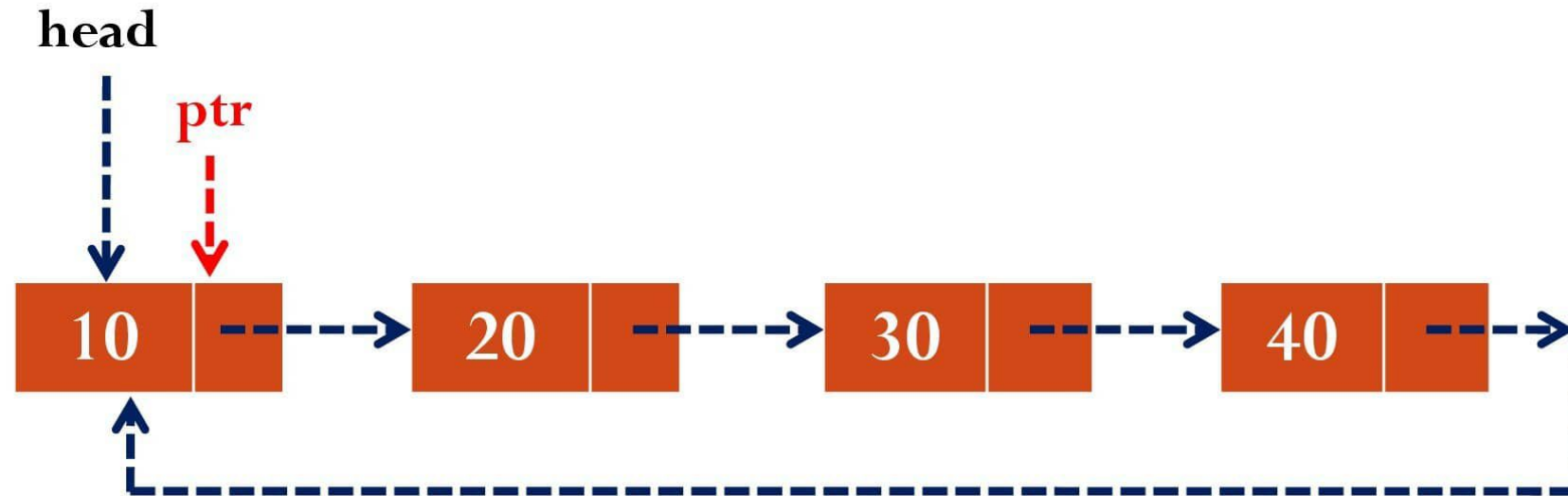
- **Disadvantages**

- When we are unable to detect the end of the list while moving from one node to the next, system may get trap into in **infinite loop**

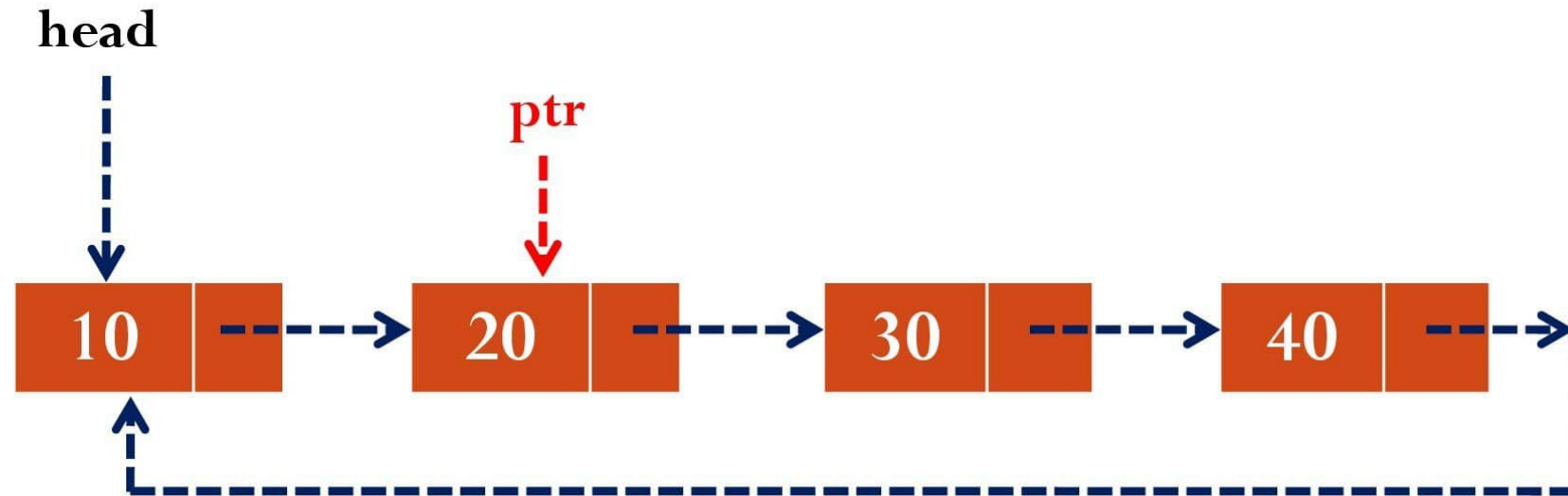
Operations on Circular Linked List

- Traverse/Display a list
- Insertion of a node into list
 - Insert at front
 - Insert at end
 - Insert after a specified node
- Deletion of node from list
 - Delete from front
 - Delete from end
 - Delete a specified node
- Searching for an element in a list
- Merging two linked list into larger list

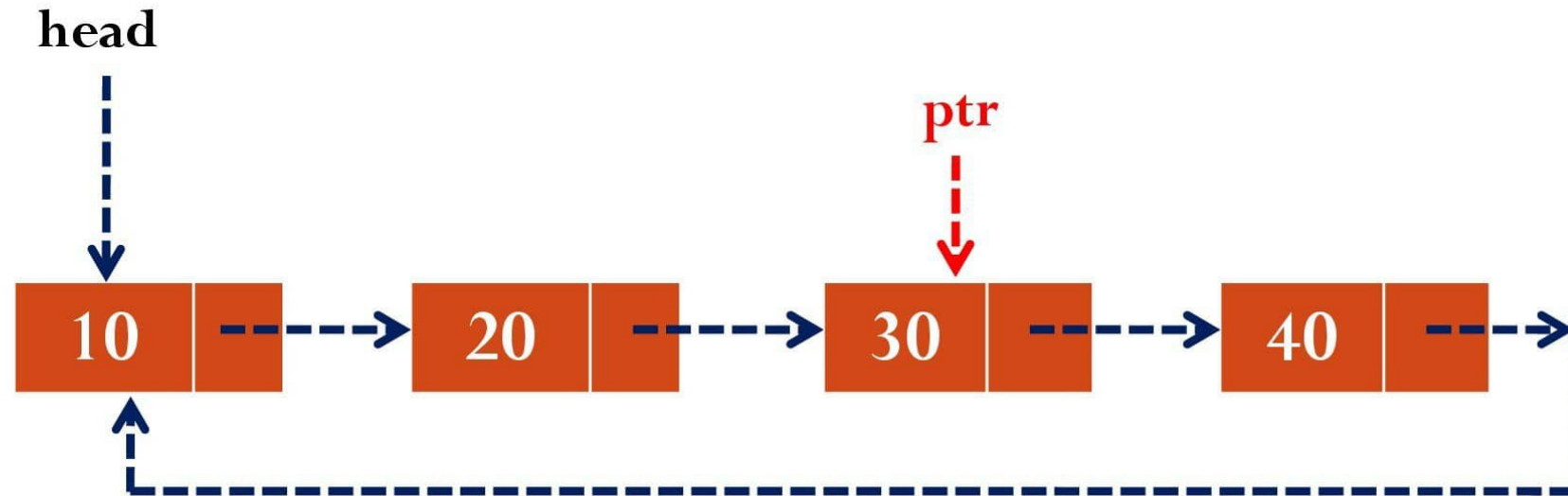
Display/Traversal



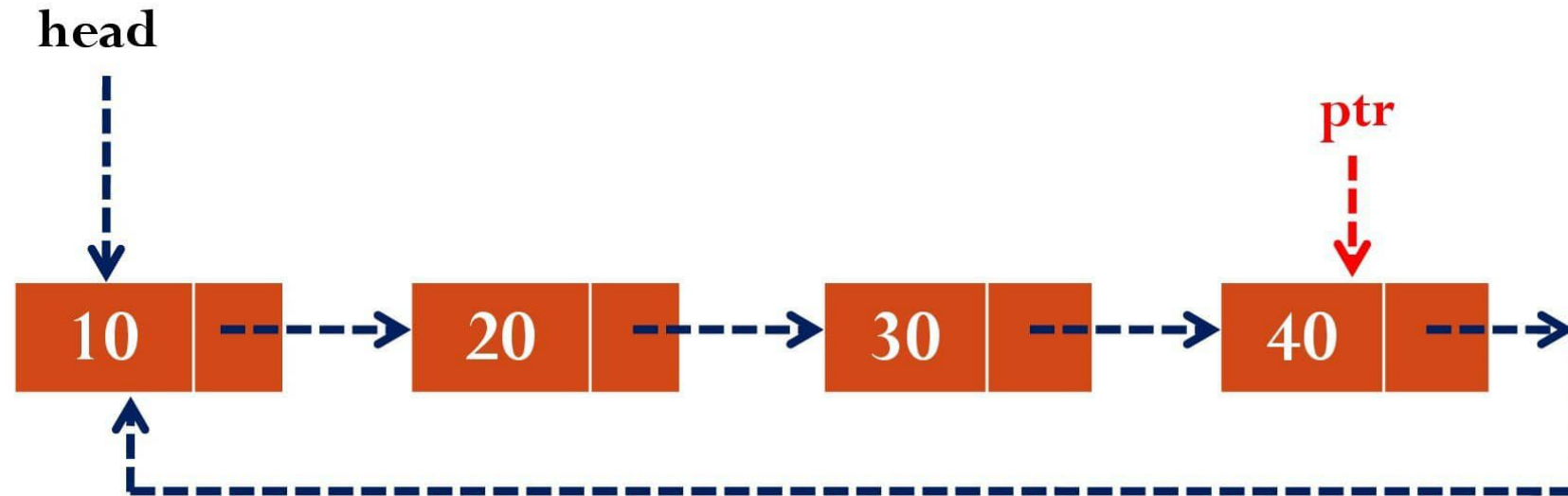
Display/Traversal



Display/Traversal



Display/Traversal



Display/Traversal - Algorithm

Algorithm Display(head)

1. If head=NULL then
 1. Print “List is Empty”
2. Else
 1. ptr=head
 2. While ptr→link≠head do
 1. Print ptr→data
 2. ptr=ptr→link
 3. Print ptr→data

Insertion

1. Insert at Front
2. Insert at End
3. Insert after a specified node

Insertion

1. Insert at Front
2. Insert at End
3. Insert after a specified node

Insert at Front

2 cases:

1. List is empty
2. List is not empty

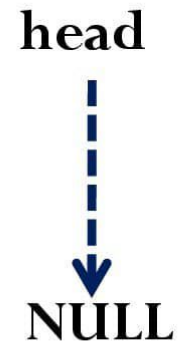
Case 1

Insert at Front



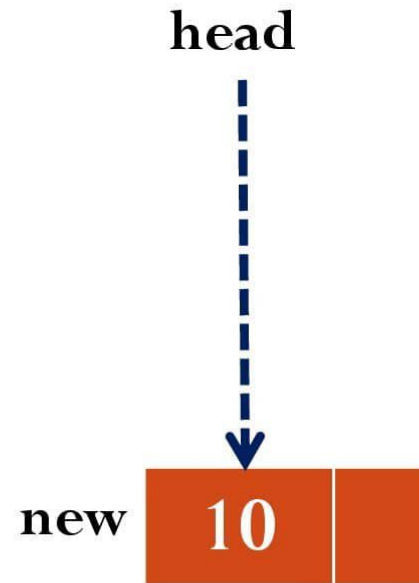
Case 1

Insert at Front



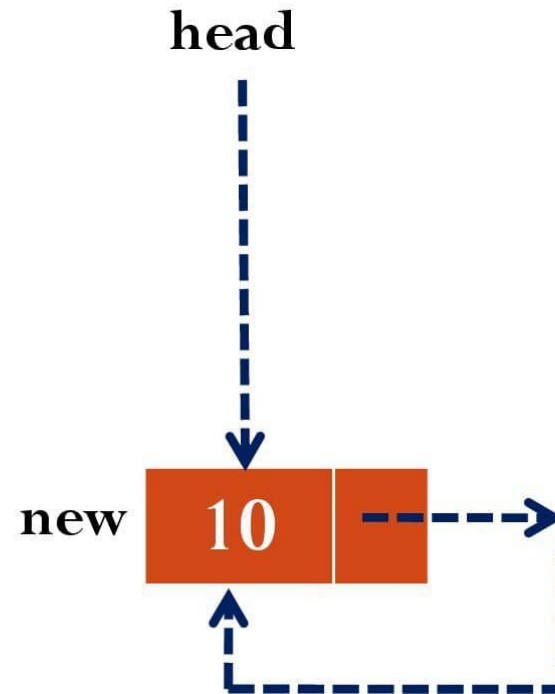
Case 1

Insert at Front



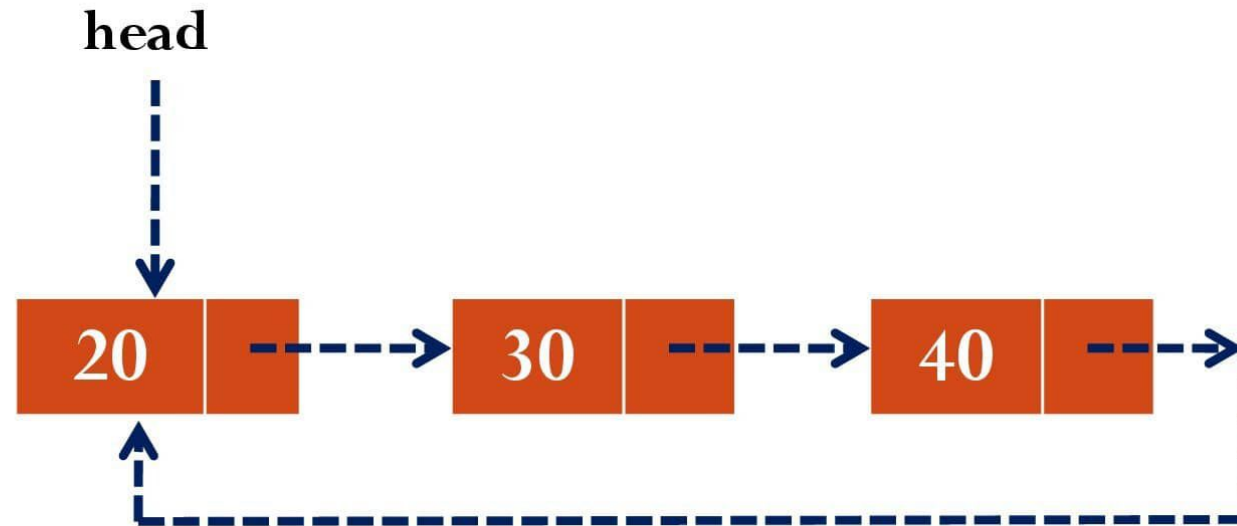
Case 1

Insert at Front



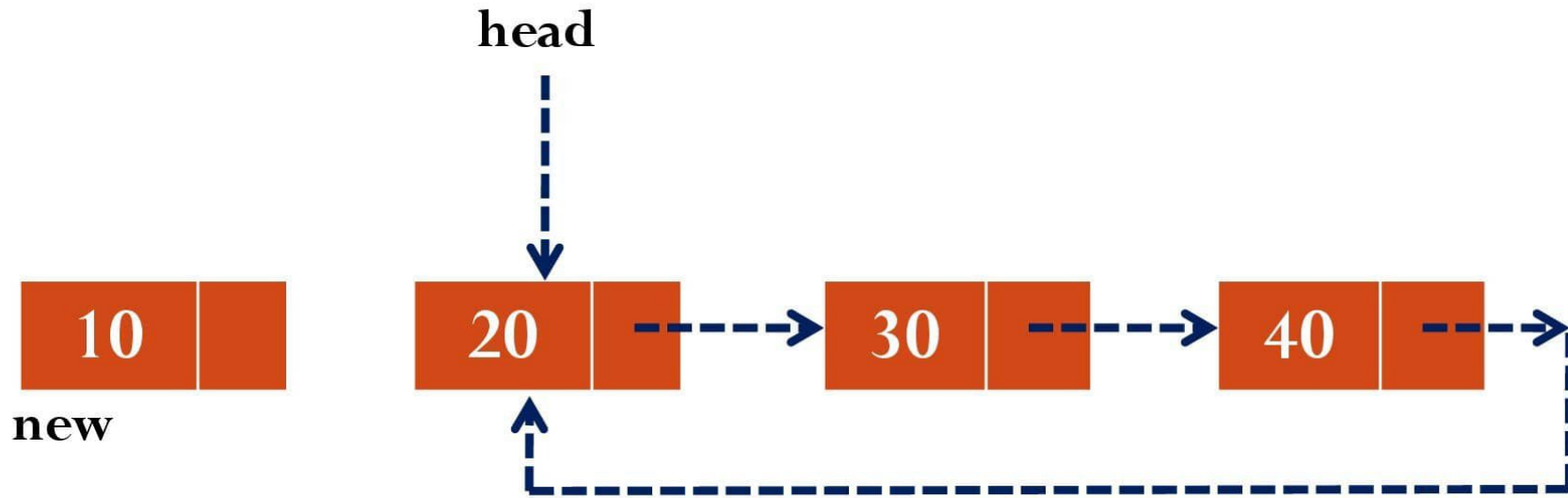
Case 2

Insert at Front



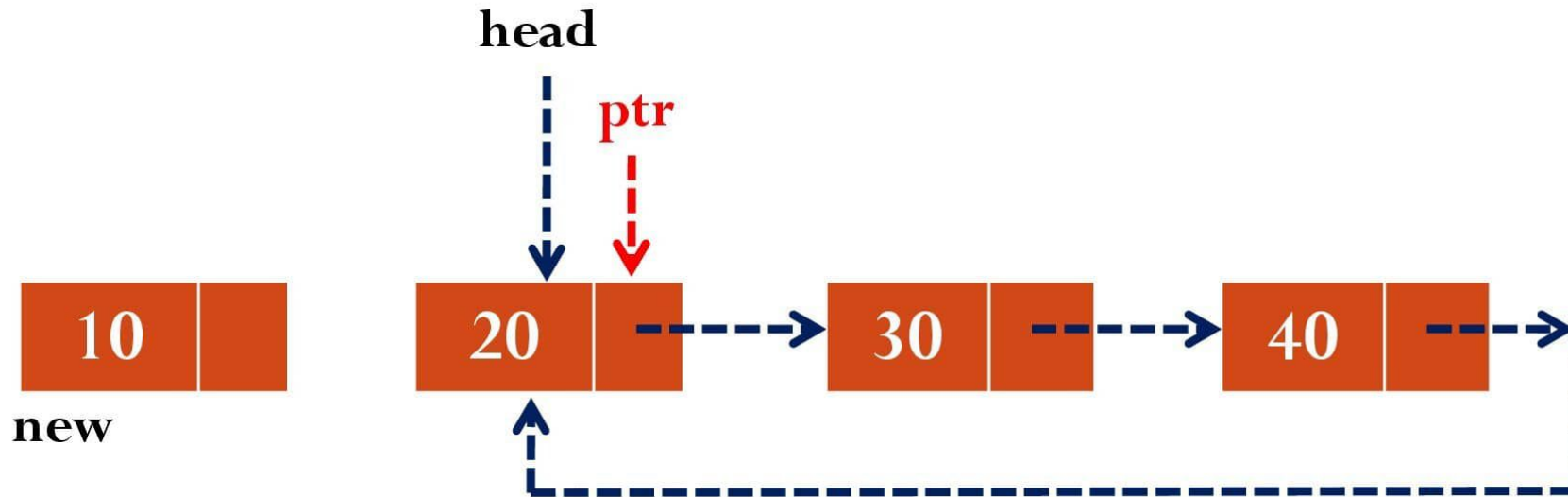
Case 2

Insert at Front



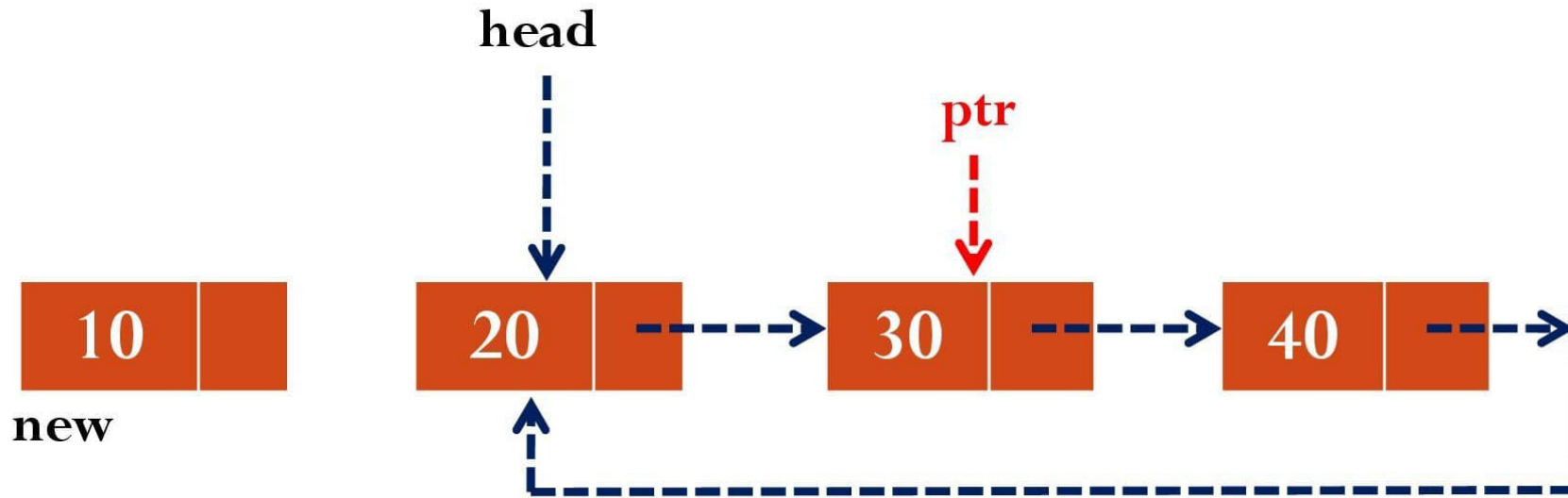
Case 2

Insert at Front



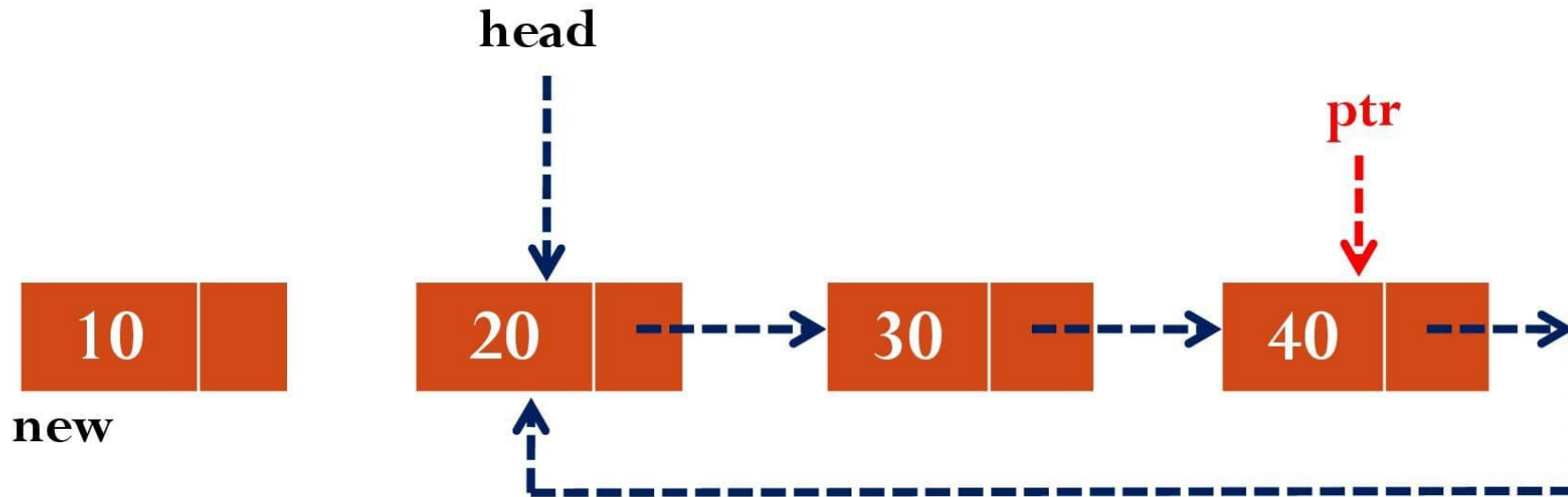
Case 2

Insert at Front



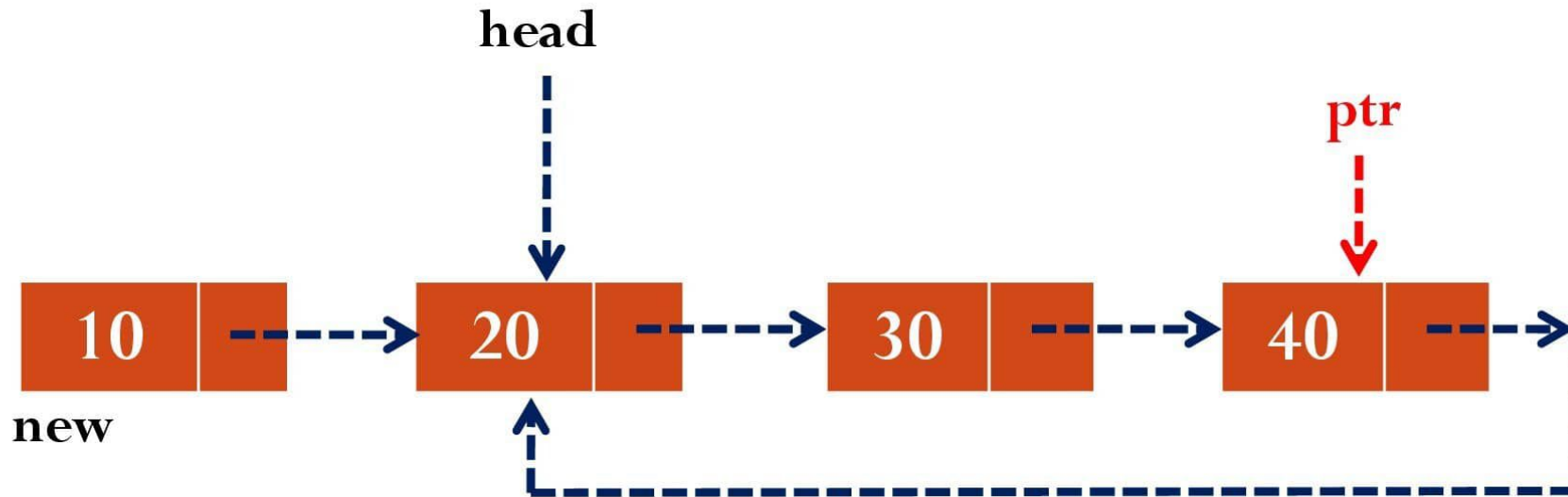
Case 2

Insert at Front



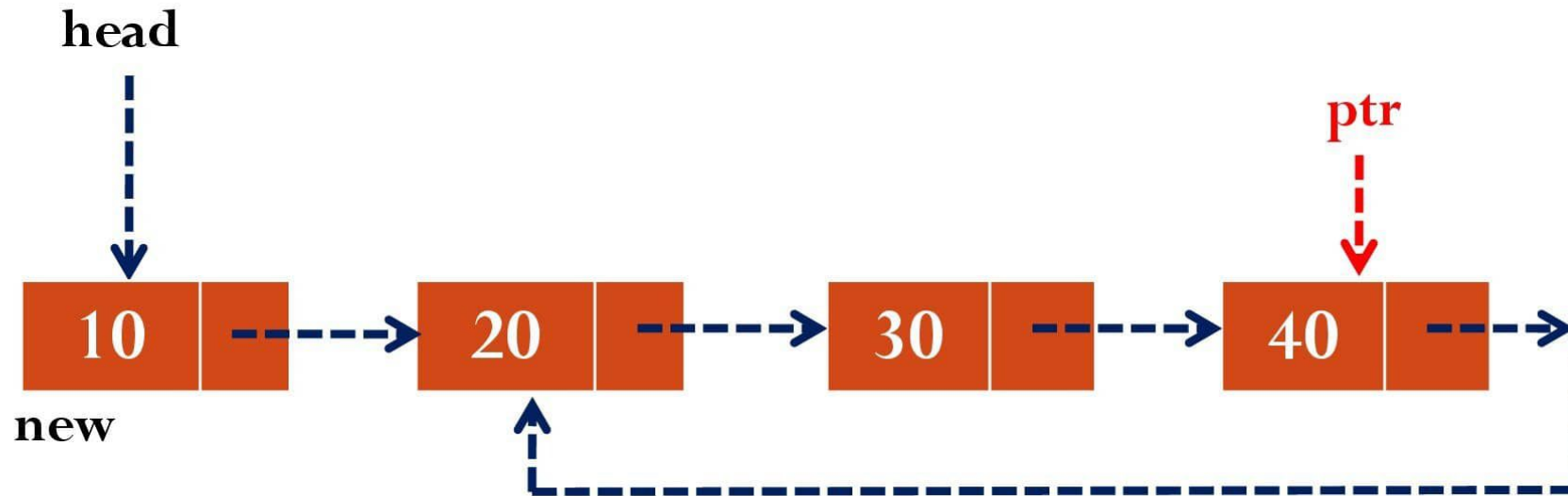
Case 2

Insert at Front



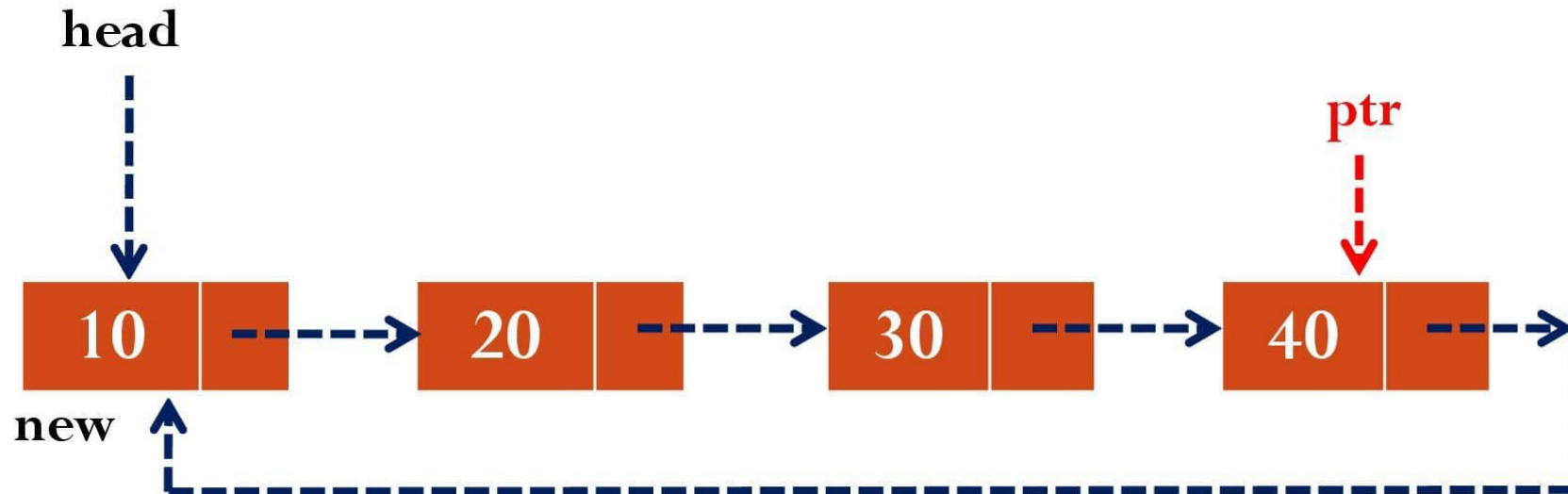
Case 2

Insert at Front



Case 2

Insert at Front



Insert at Front ~ Algorithm

Algorithm Insert_Front(head, x)

1. Create a node new
2. $\text{new} \rightarrow \text{data} = x$
3. If $\text{head} = \text{NULL}$ then
 1. $\text{head} = \text{new}$
 2. $\text{new} \rightarrow \text{link} = \text{new}$
4. Else
 1. $\text{ptr} = \text{head}$
 2. While $\text{ptr} \rightarrow \text{link} \neq \text{head}$ do
 1. $\text{ptr} = \text{ptr} \rightarrow \text{link}$
 3. $\text{new} \rightarrow \text{link} = \text{head}$
 4. $\text{head} = \text{new}$
 5. $\text{ptr} \rightarrow \text{link} = \text{head}$

Insertion

1. Insert at Front
2. **Insert at End**
3. Insert after a specified node

Insert at End

2 cases:

1. List is empty
2. List is not empty

Case 1

Insert at End



Case 1

Insert at End

head



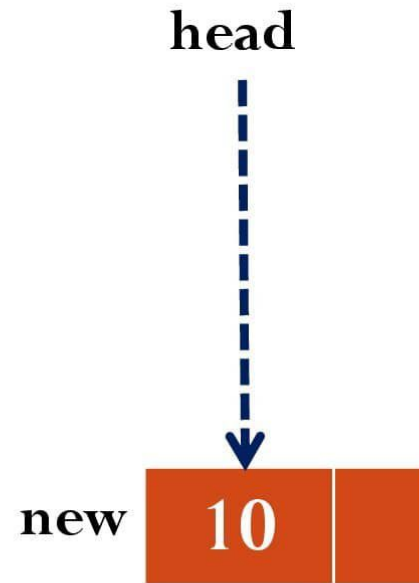
NULL

new

10

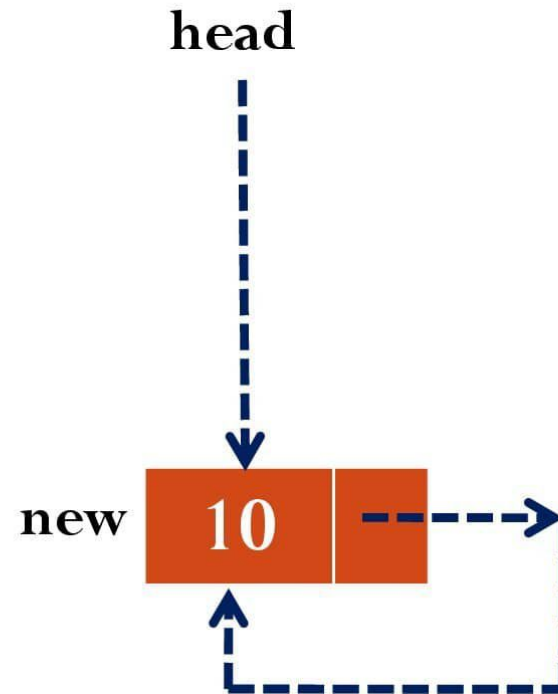
Case 1

Insert at End



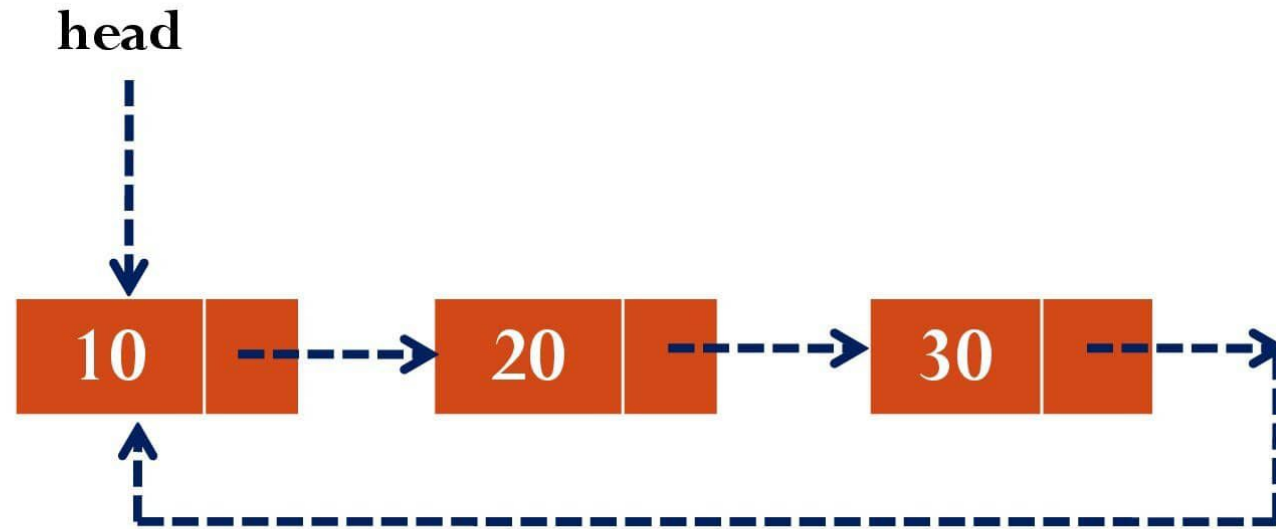
Case 1

Insert at End



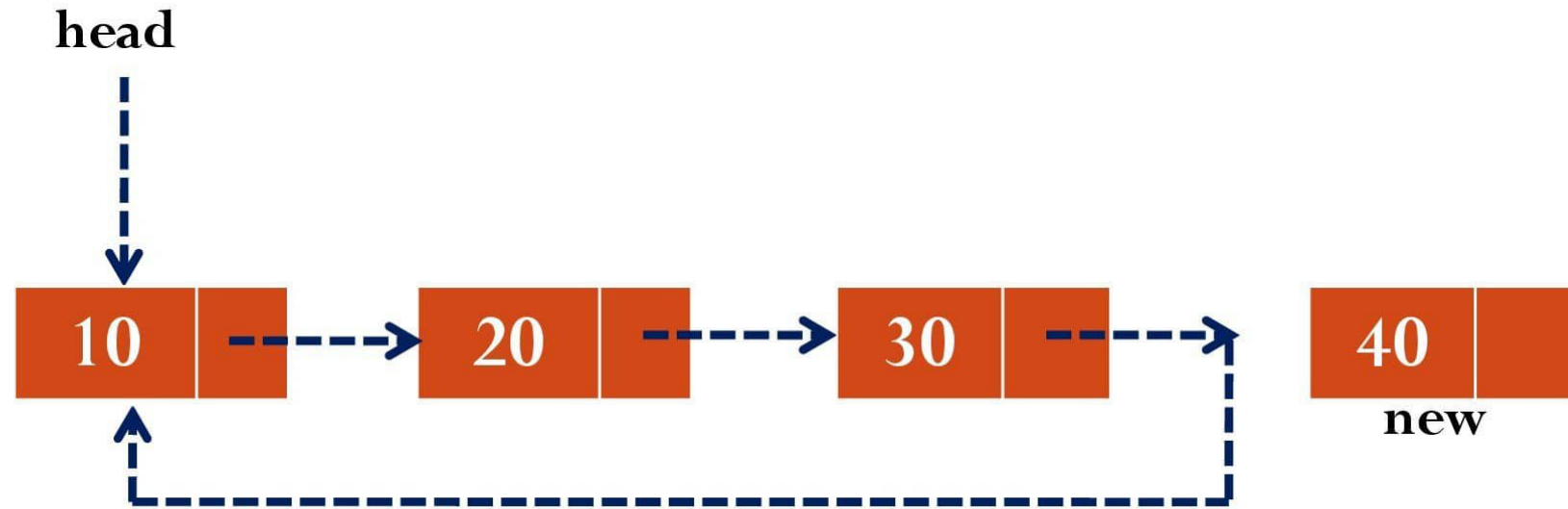
Case 2

Insert at End



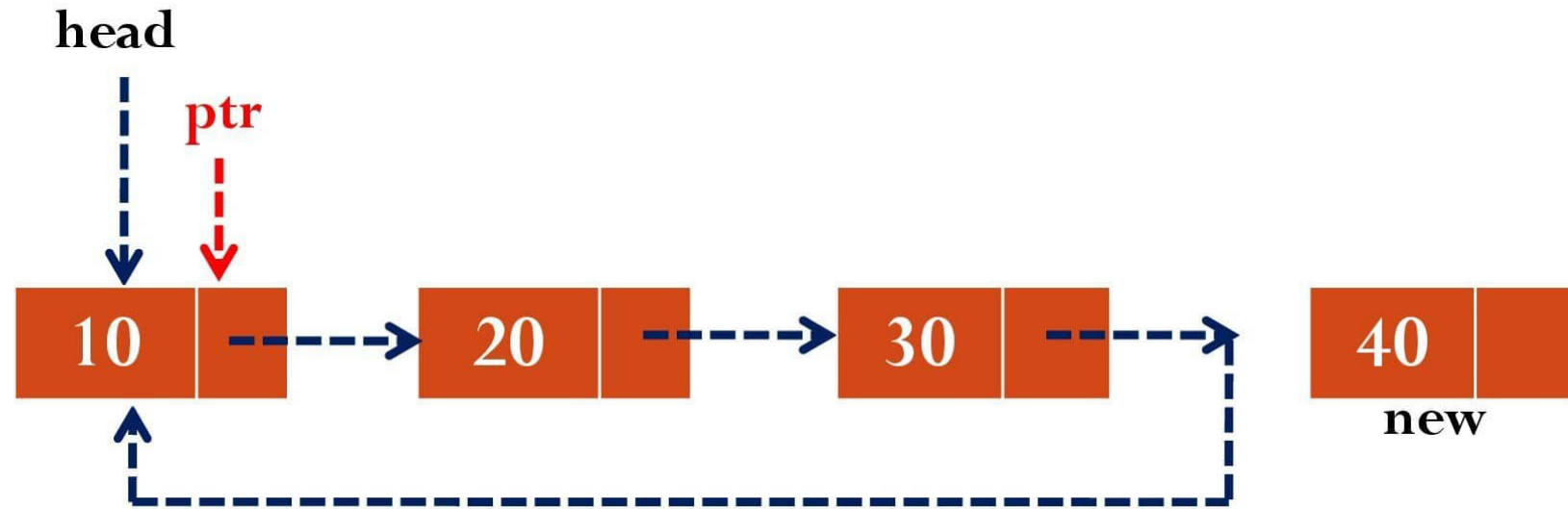
Case 2

Insert at End



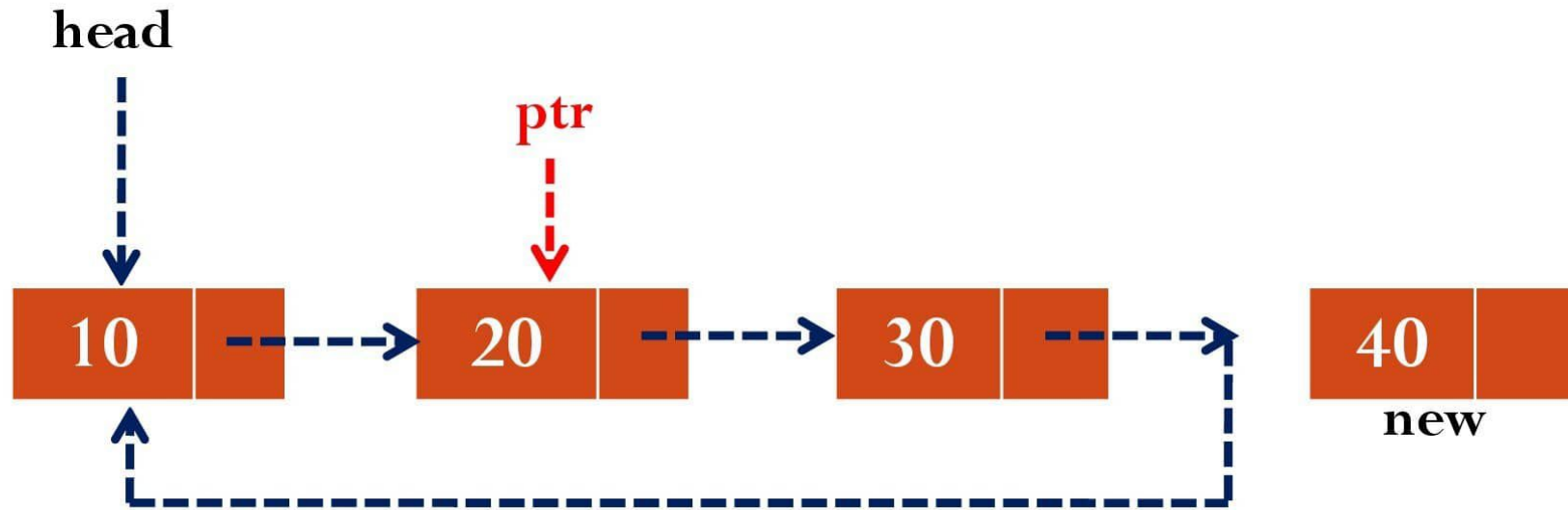
Case 2

Insert at End



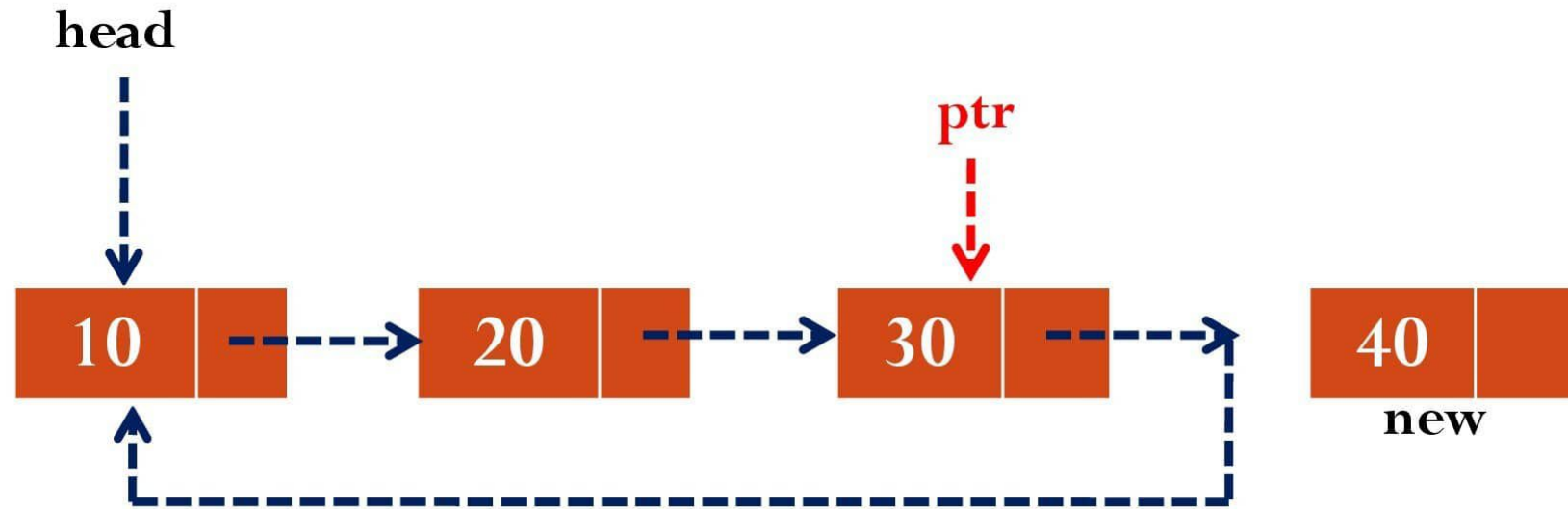
Case 2

Insert at End



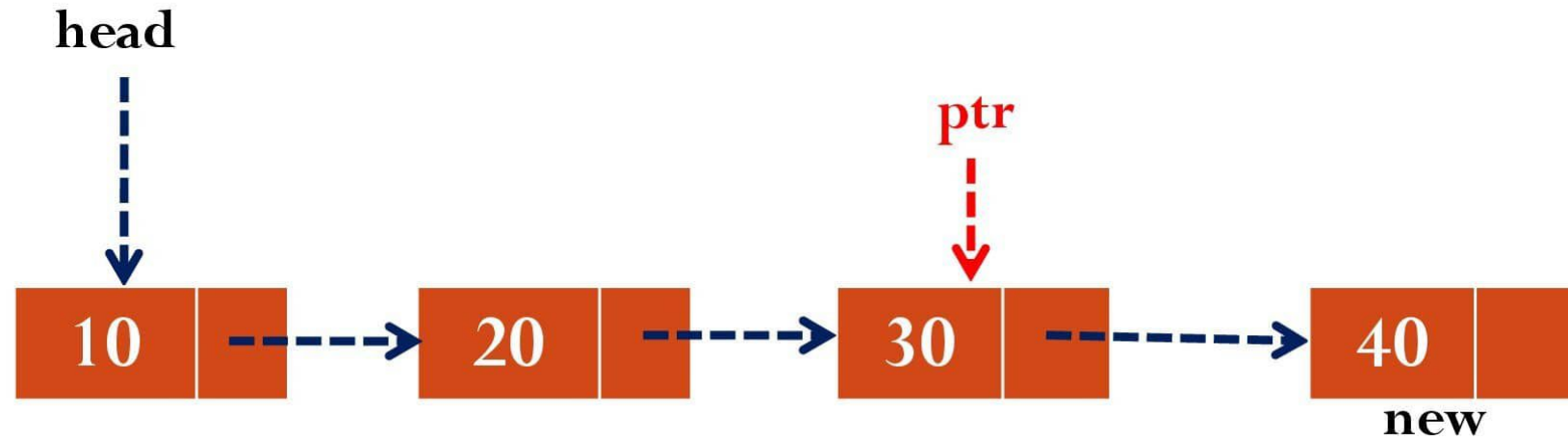
Case 2

Insert at End



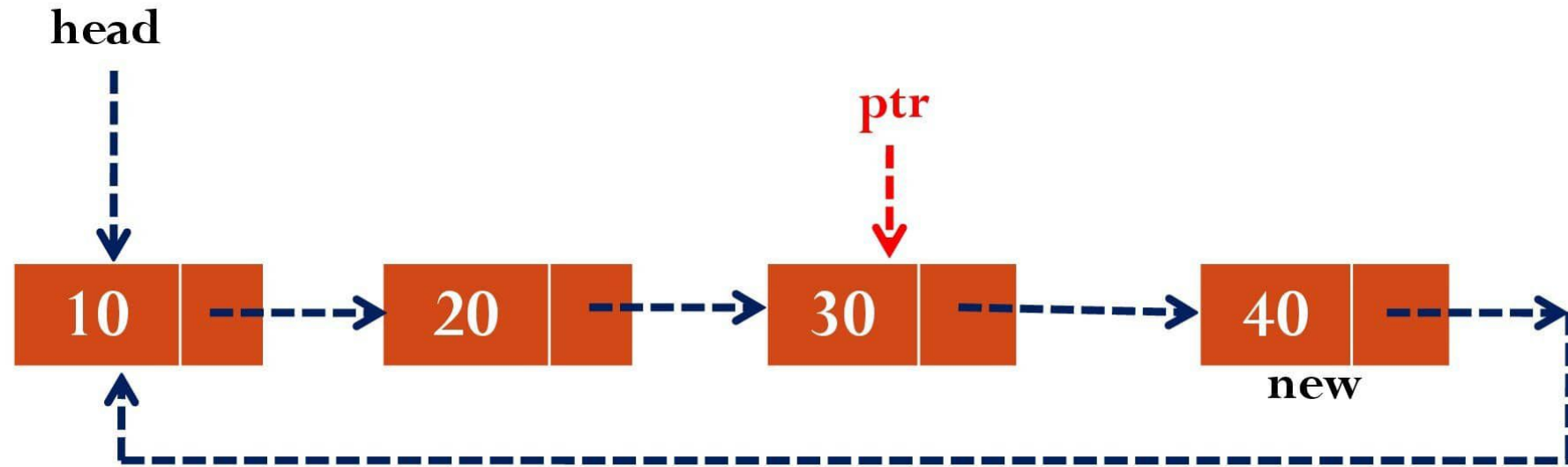
Case 2

Insert at End



Case 2

Insert at End



Insert at End- Algorithm

Algorithm Insert_End(head, x)

1. Create a node new
2. $\text{new} \rightarrow \text{data} = x$
3. If $\text{head} = \text{NULL}$ then
 1. $\text{new} \rightarrow \text{link} = \text{new}$
 2. $\text{head} = \text{new}$
4. Else
 1. $\text{ptr} = \text{head}$
 2. While $\text{ptr} \rightarrow \text{link} \neq \text{head}$ do
 1. $\text{ptr} = \text{ptr} \rightarrow \text{link}$
 3. $\text{ptr} \rightarrow \text{link} = \text{new}$
 4. $\text{new} \rightarrow \text{link} = \text{head}$

Insertion

1. Insert at Front
2. Insert at End
3. **Insert after a specified node**

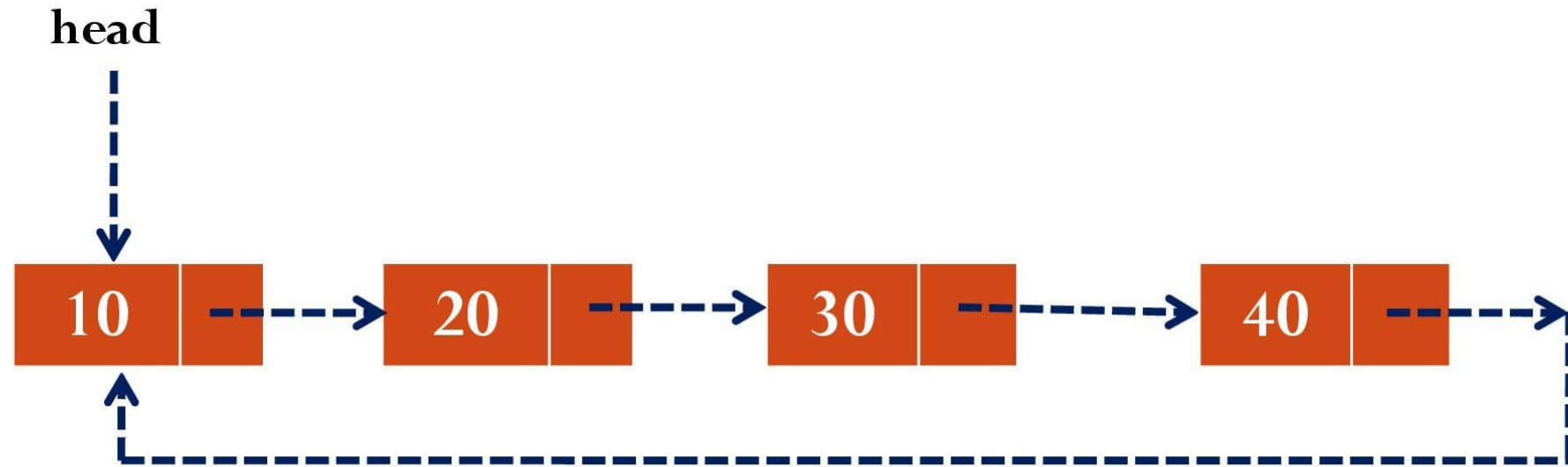
Insert after a specified node

2 cases:

1. List is empty
2. List is not empty

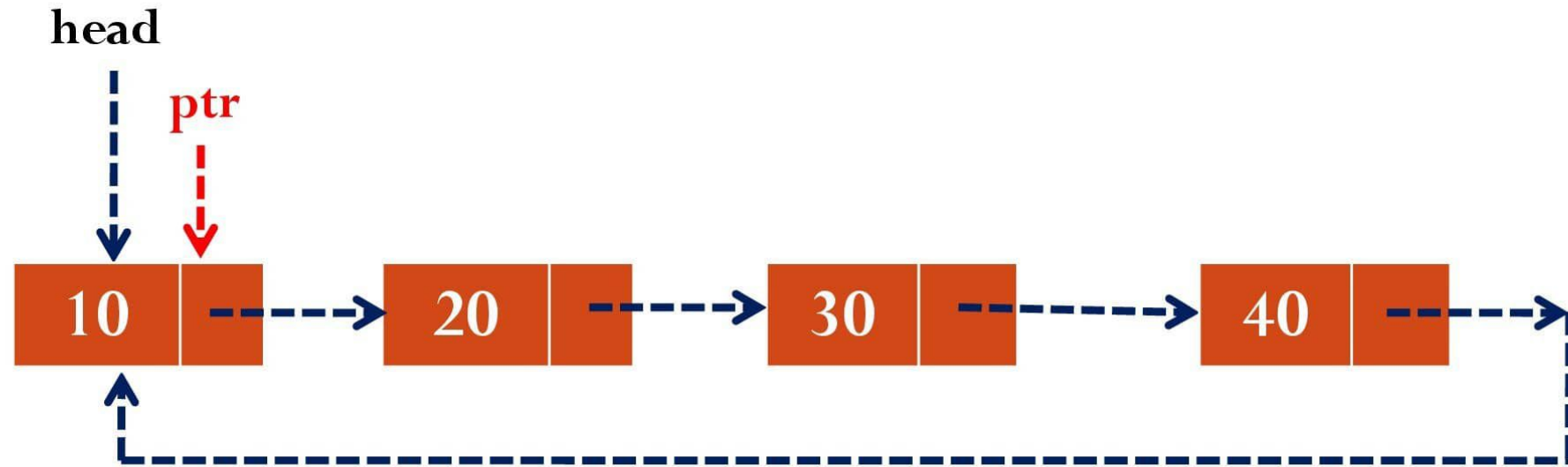
Insert after 30

Case 2



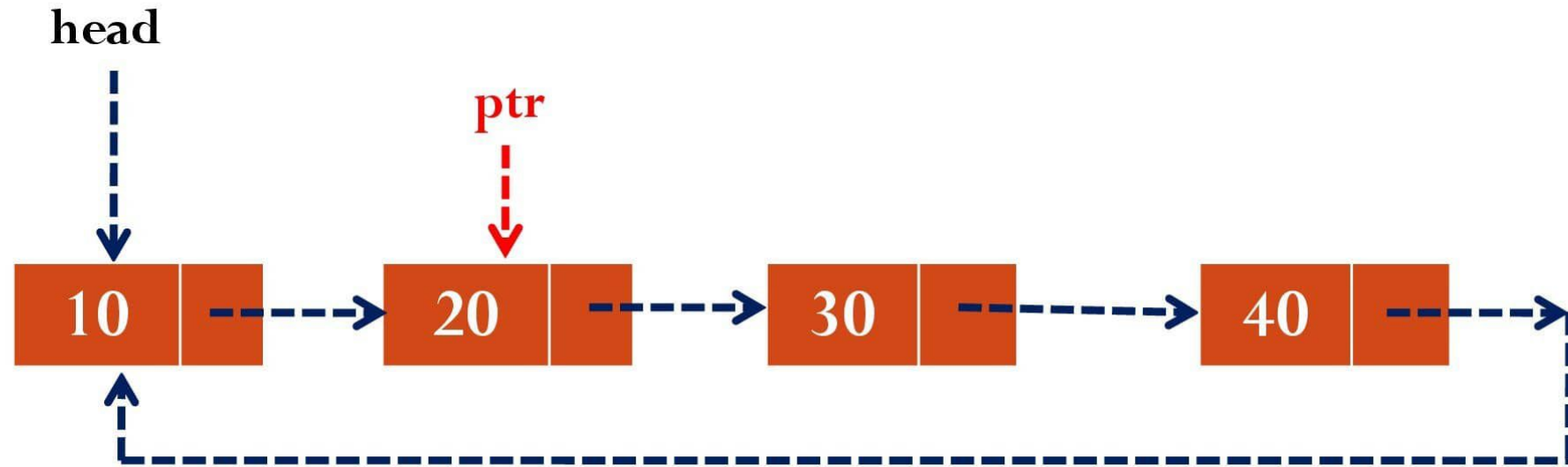
Insert after 30

Case 2



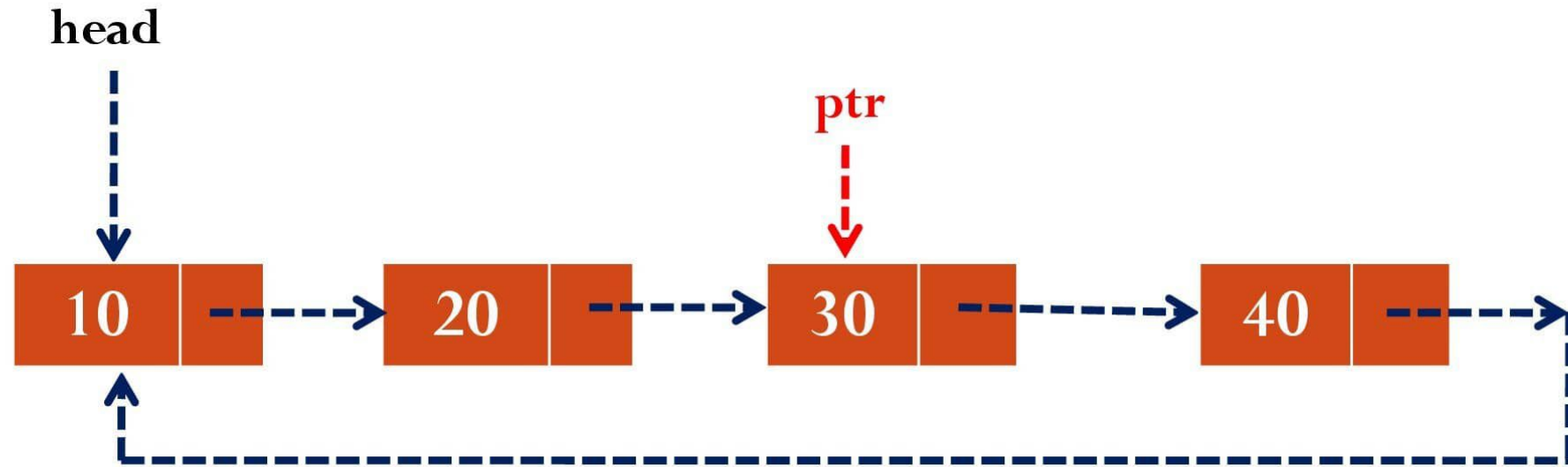
Insert after 30

Case 2



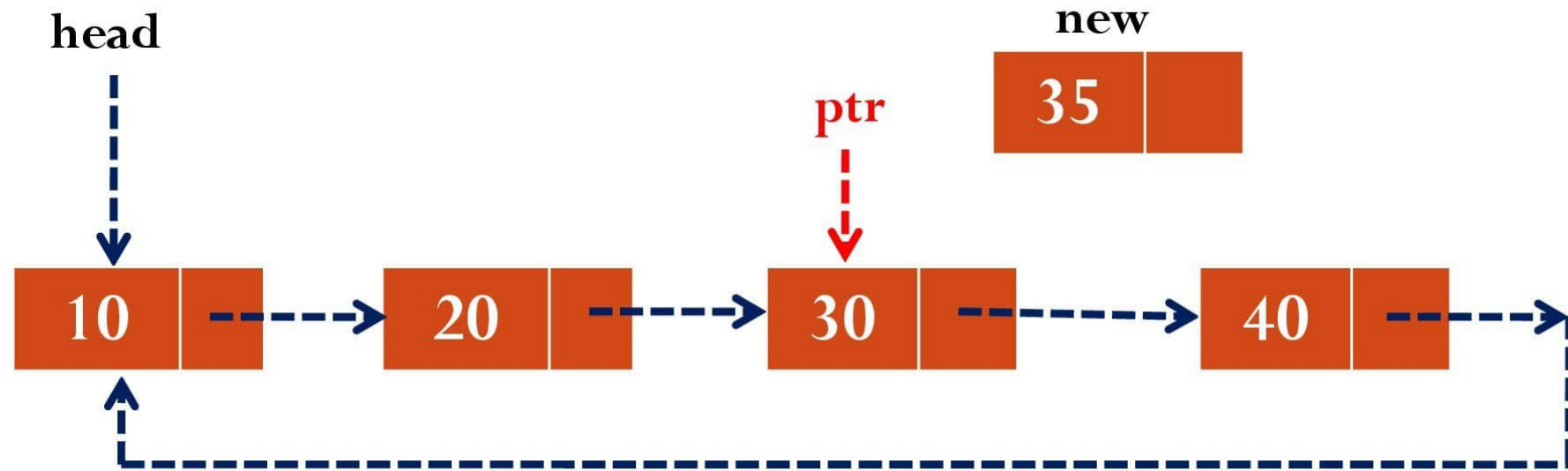
Insert after 30

Case 2



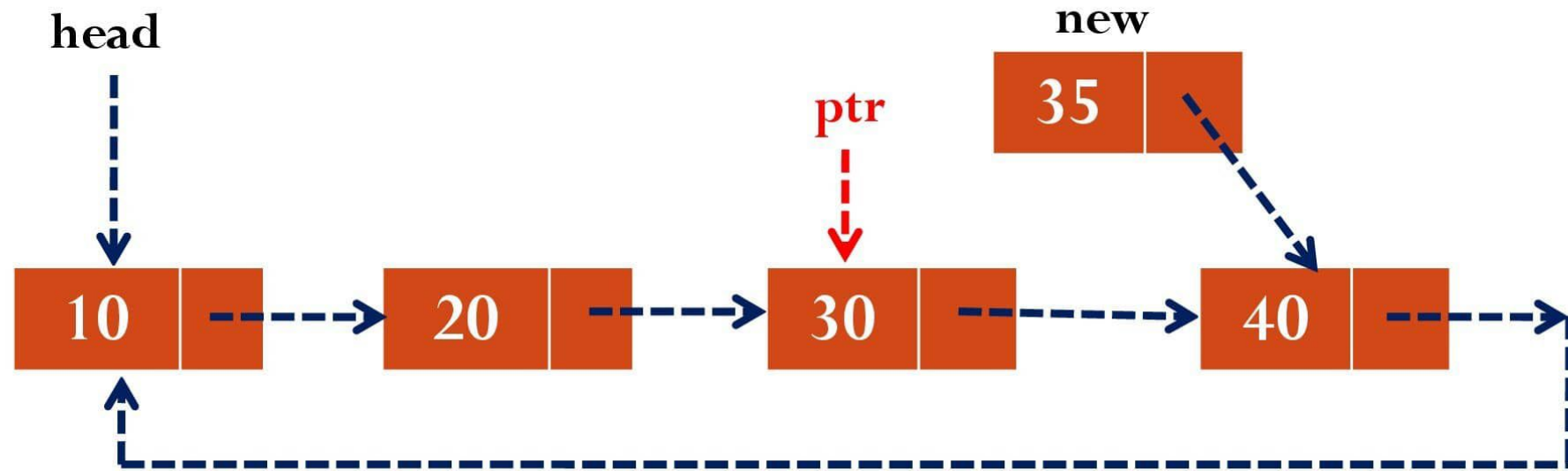
Insert after 30

Case 2



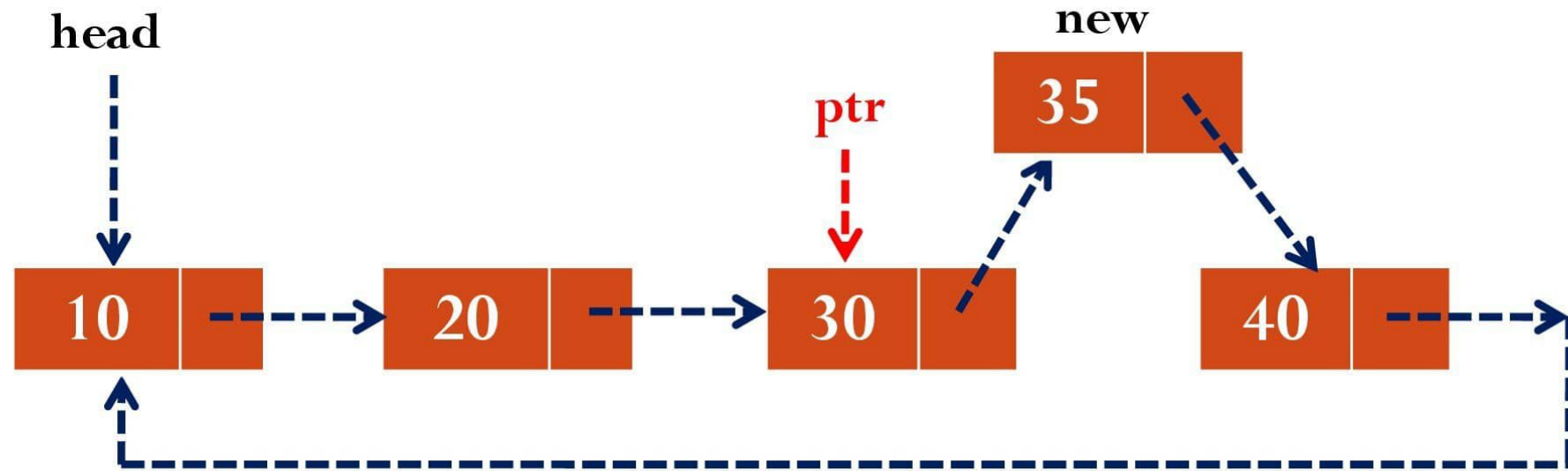
Insert after 30

Case 2



Insert after 30

Case 2



Insert after a specified node ~ Algorithm

Algorithm Insert_After(head, key,x)

1. If head=NULL then
 1. Print “Search data not found”
2. Else
 1. ptr=head
 2. While ptr→data!=key do
 1. ptr=ptr→link
 2. If ptr=head then
 1. Break
3. If ptr→data=key then
 1. Create a node new
 2. new→data=x
 3. new→link=ptr→link
 4. ptr→link=new
4. Else
 3. Print “Search data not found. Insertion not possible.”

Deletion

1. Delete from Front
2. Delete from End
3. Delete a specified node

Deletion

1. Delete from Front
2. Delete from End
3. Delete a specified node

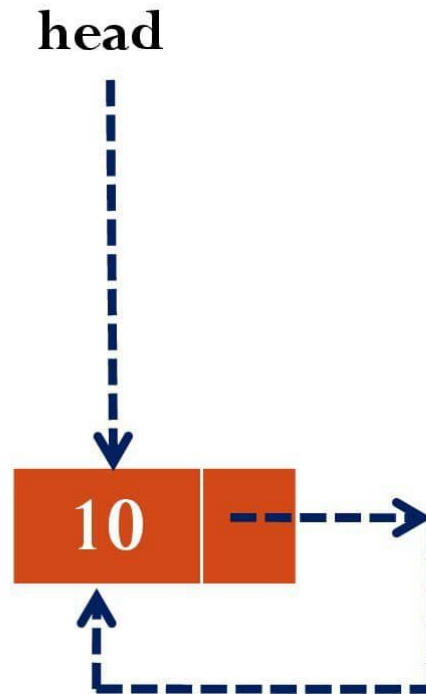
Delete from Front

2 cases:

1. List is empty
2. List contains only one node
3. List contains more than one node

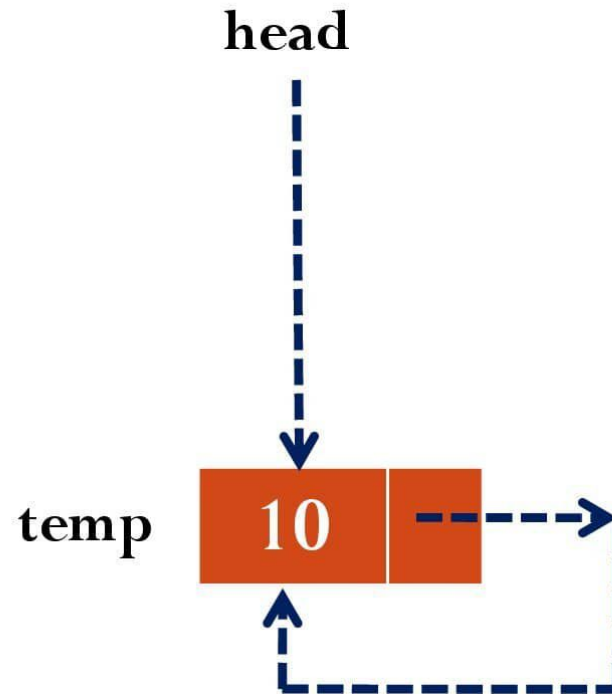
Case 2

Delete from Front



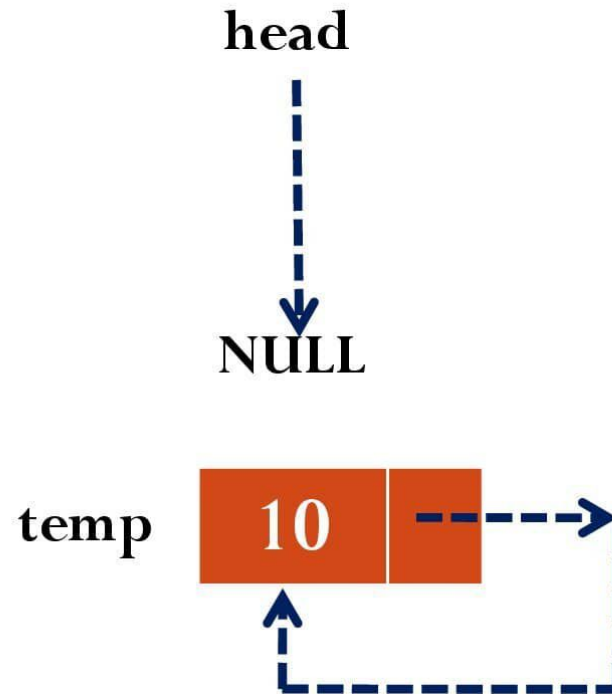
Case 2

Delete from Front



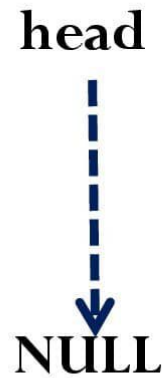
Case 2

Delete from Front



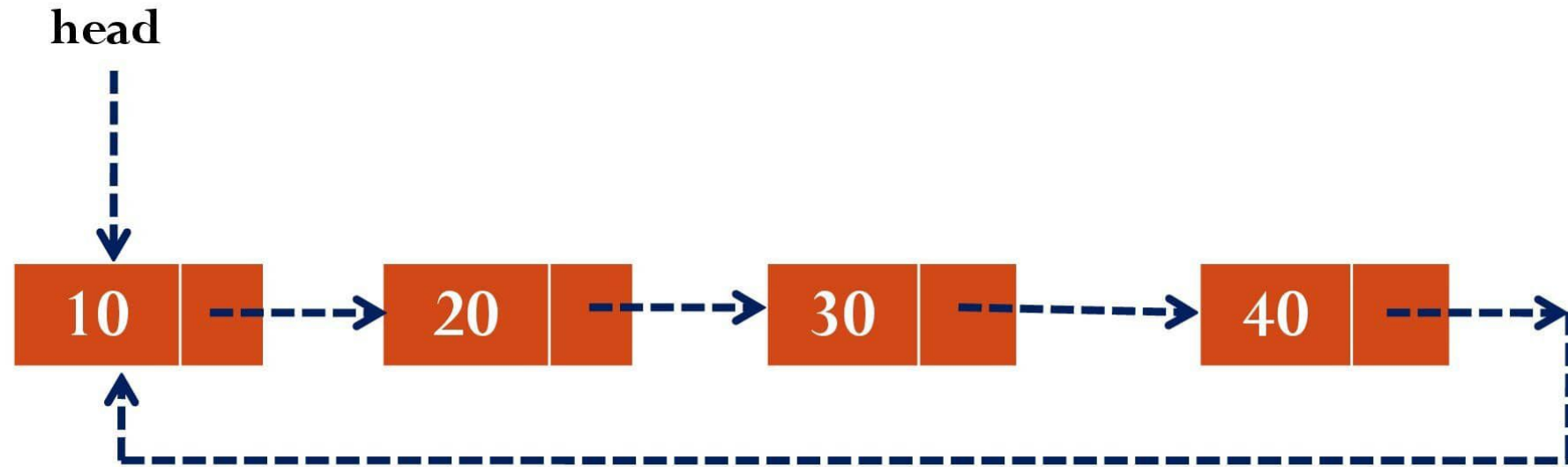
Case 2

Delete from Front



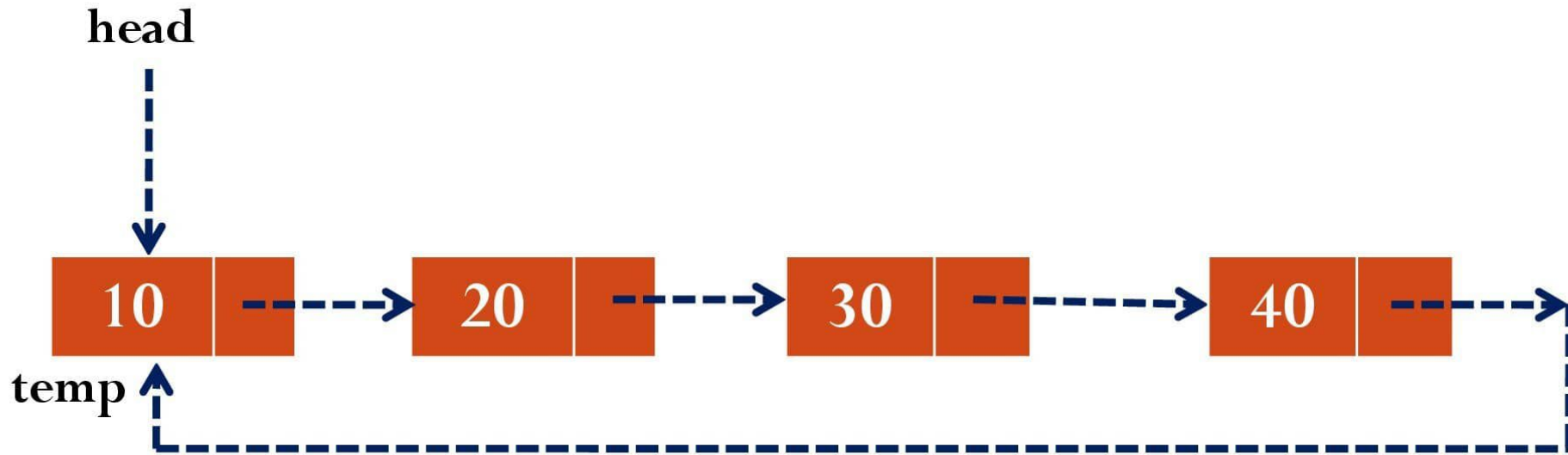
Case 3

Delete from Front



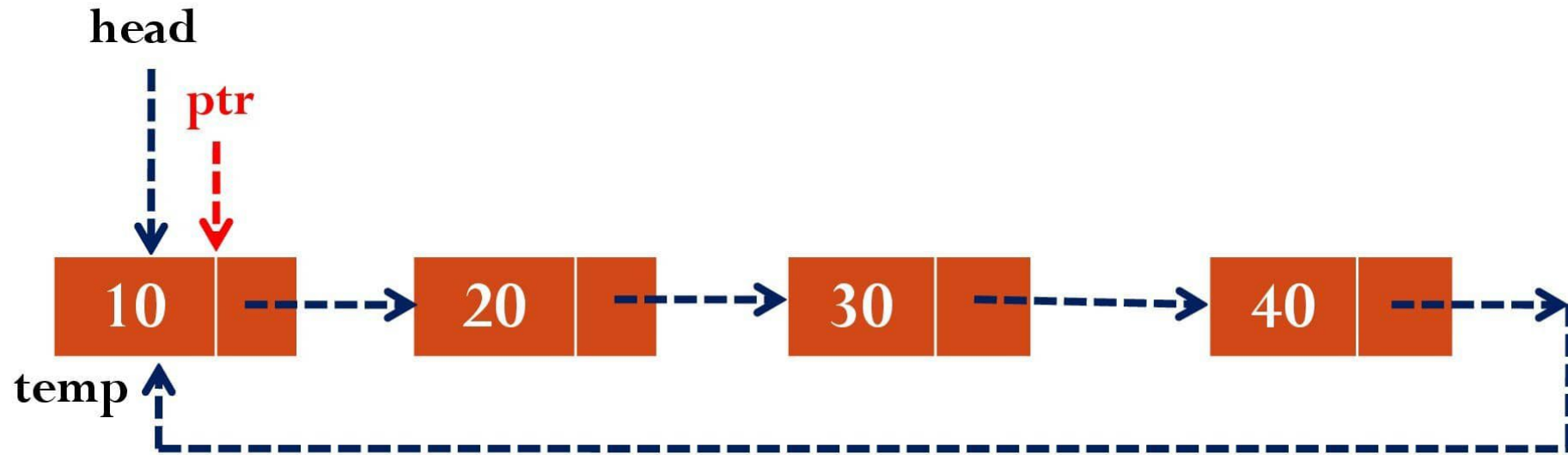
Case 3

Delete from Front



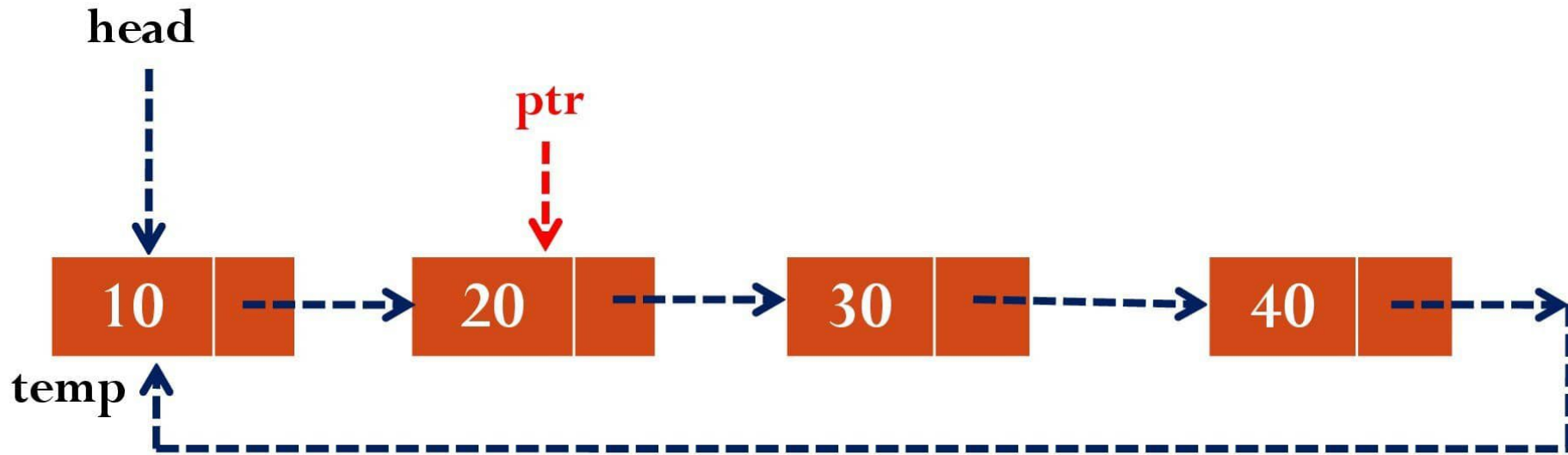
Case 3

Delete from Front



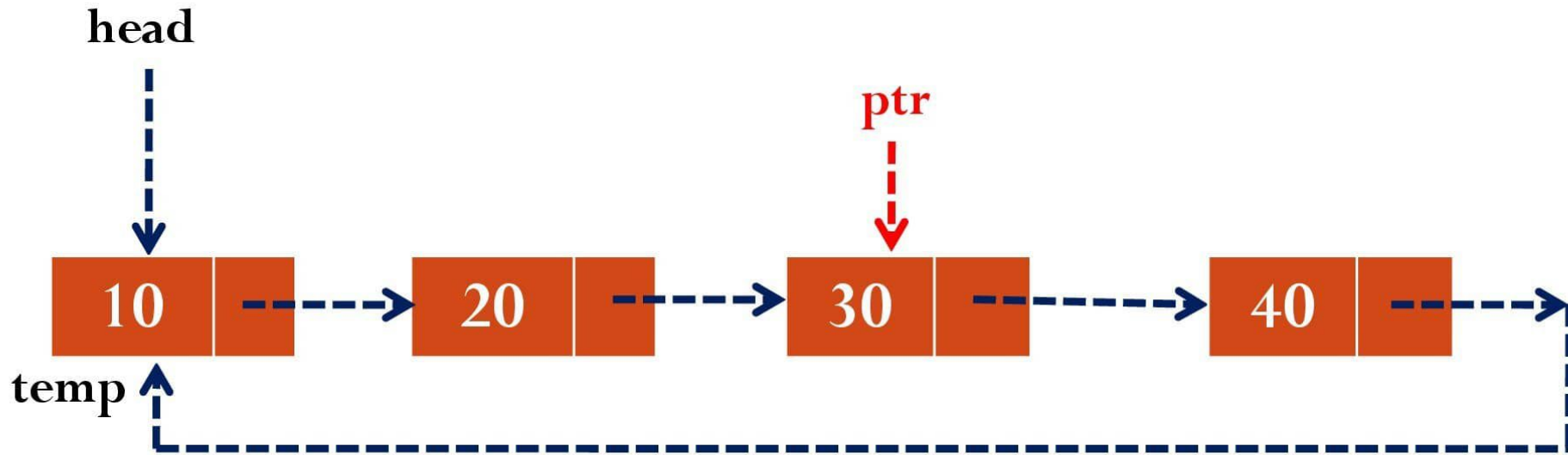
Case 3

Delete from Front



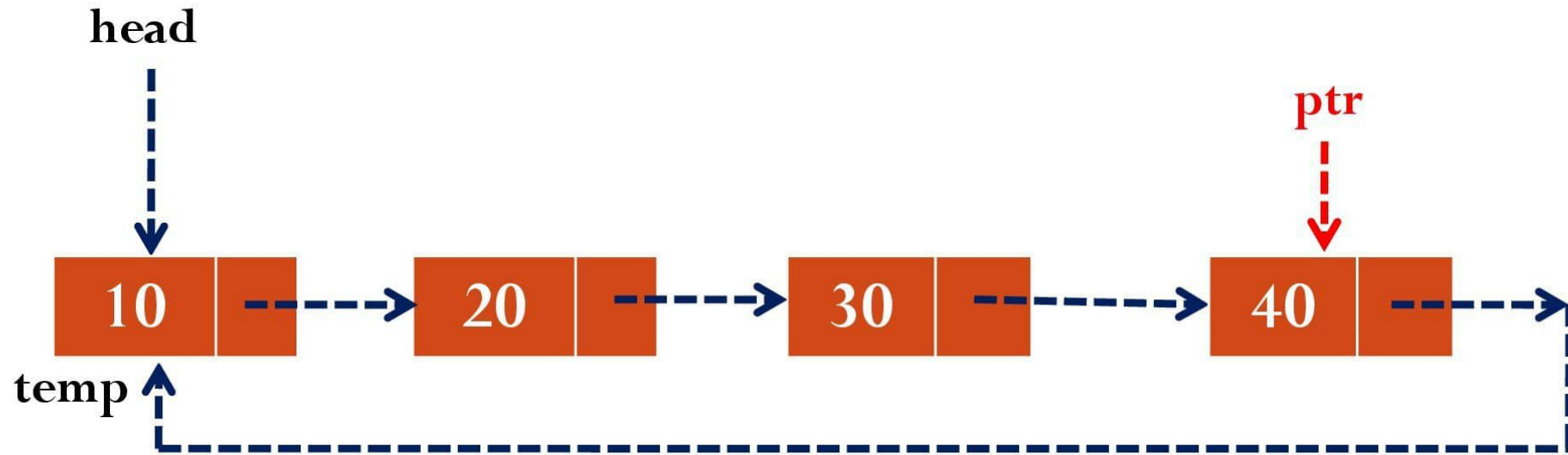
Case 3

Delete from Front



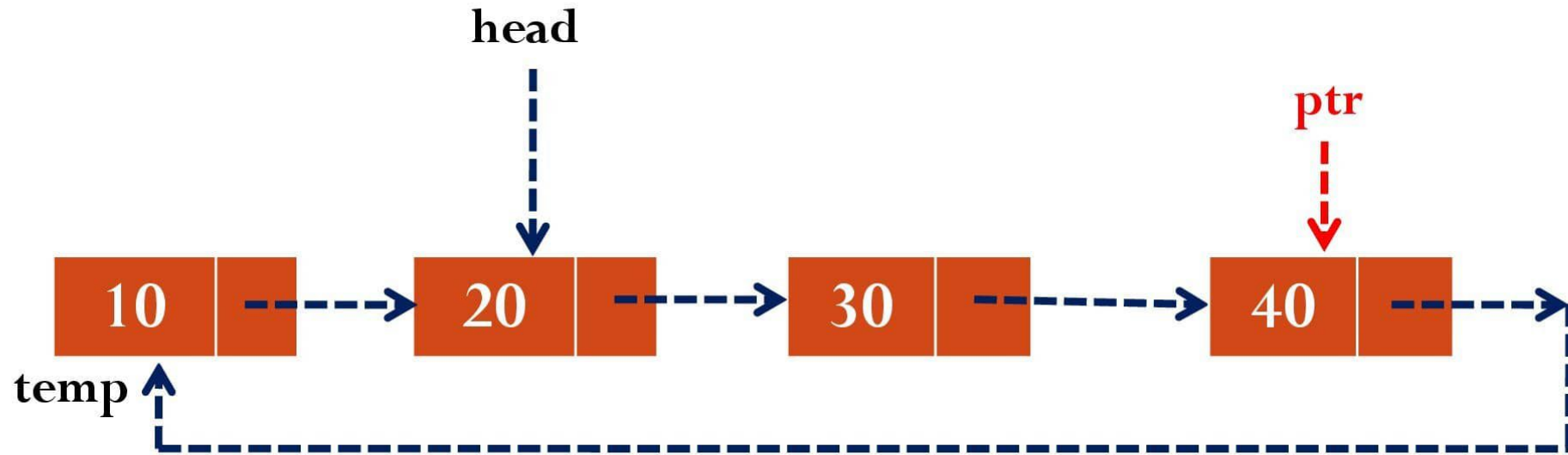
Case 3

Delete from Front



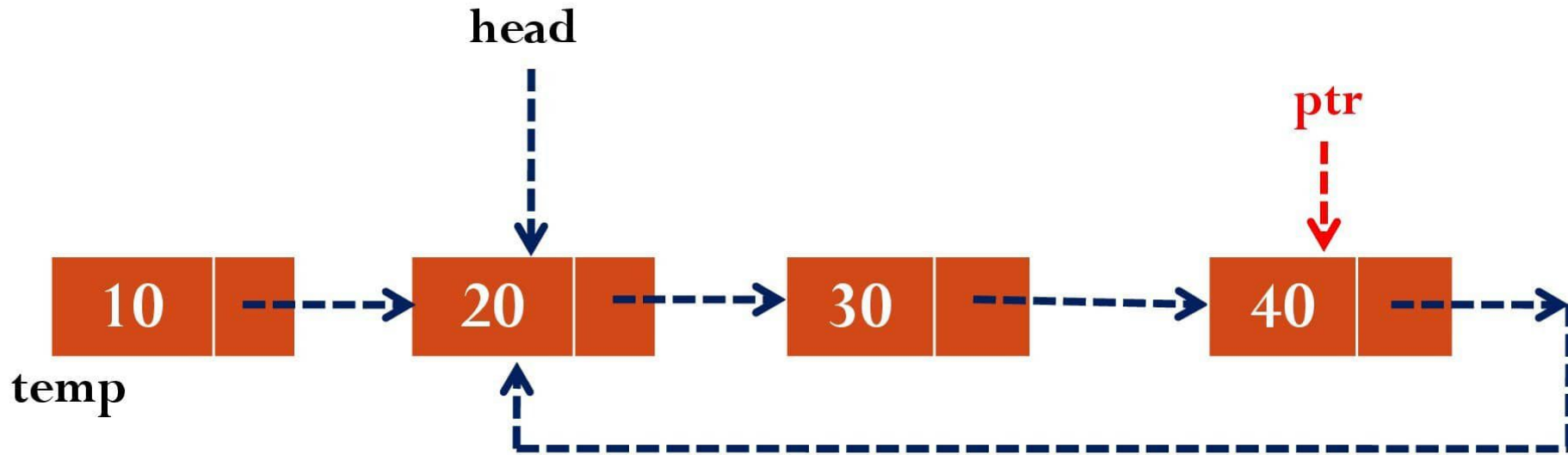
Case 3

Delete from Front



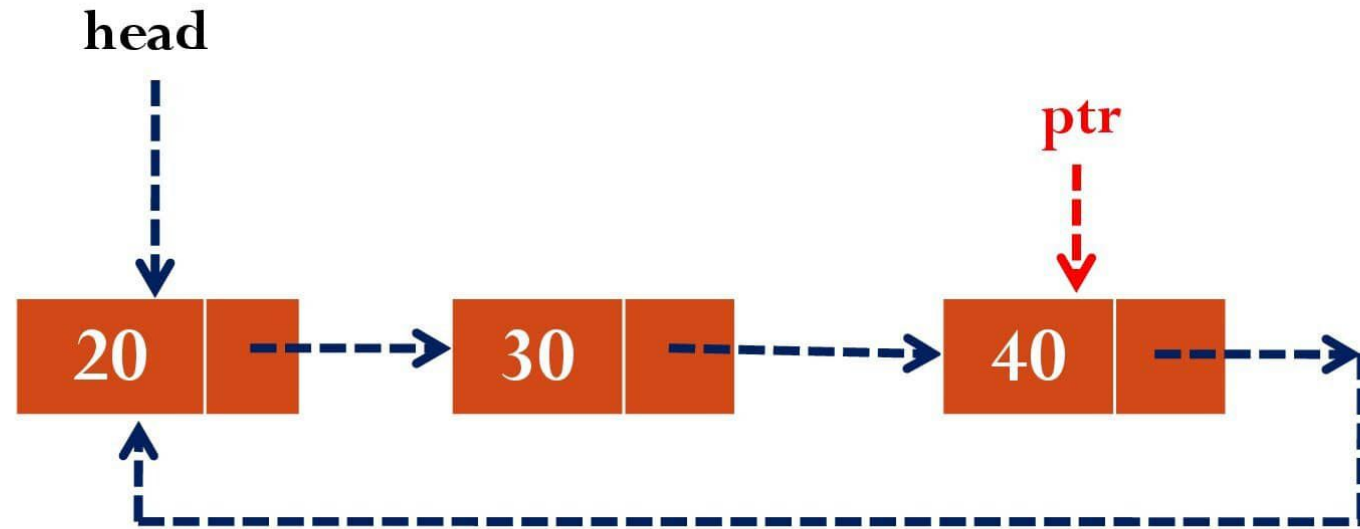
Case 3

Delete from Front



Case 3

Delete from Front



Delete from Front~ Algorithm

Algorithm Delete_Front(head)

1. If head=NULL then
 1. Print “List is Empty. Deletion is not possible”
2. Else if head→link=head then
 1. temp=head
 2. head=NULL
 3. dispose(temp)
3. Else
 1. temp=ptr=head
 2. While ptr→link≠head do
 1. ptr=ptr→link
 3. head=head→link
 4. ptr→link=head
 5. dispose(temp)

Deletion

1. Delete from Front
2. **Delete from End**
3. Delete a specified node

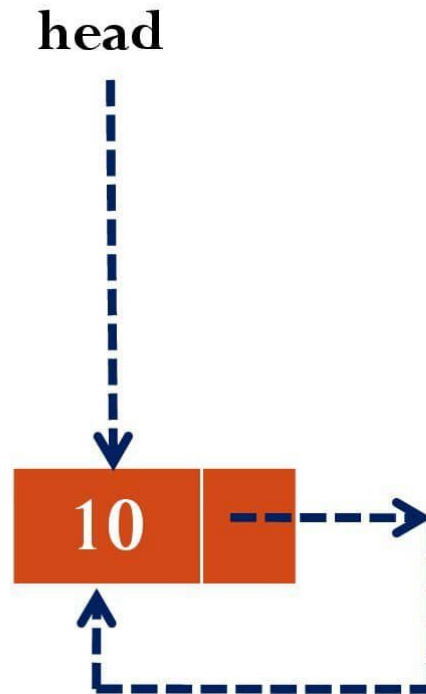
Delete from End

2 cases:

1. List is empty
2. List contains only one node
3. List contains more than one node

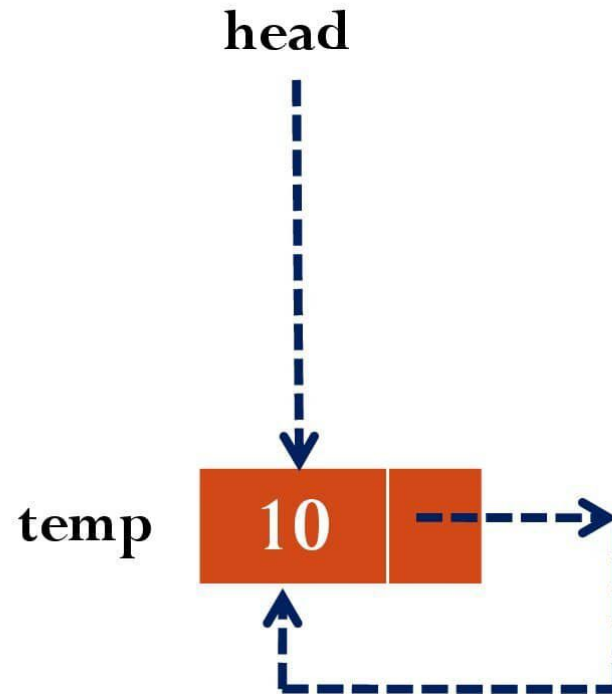
Case 2

Delete from End



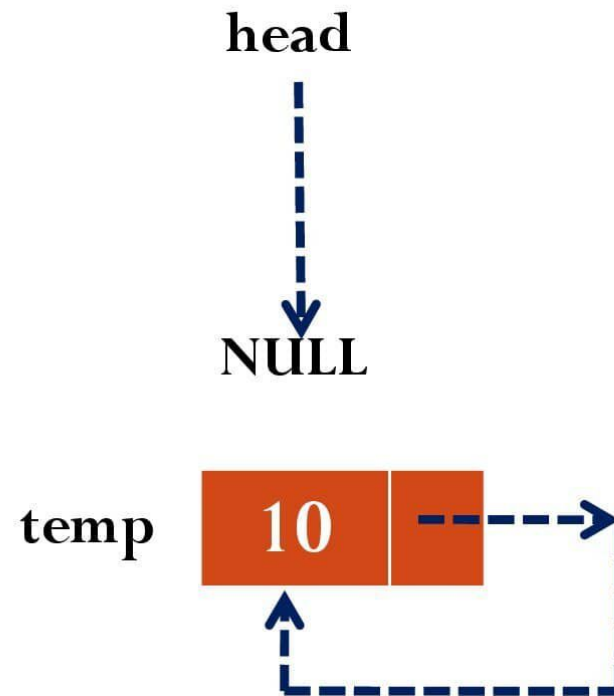
Case 2

Delete from End



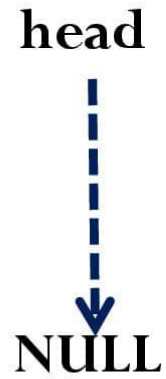
Case 2

Delete from End



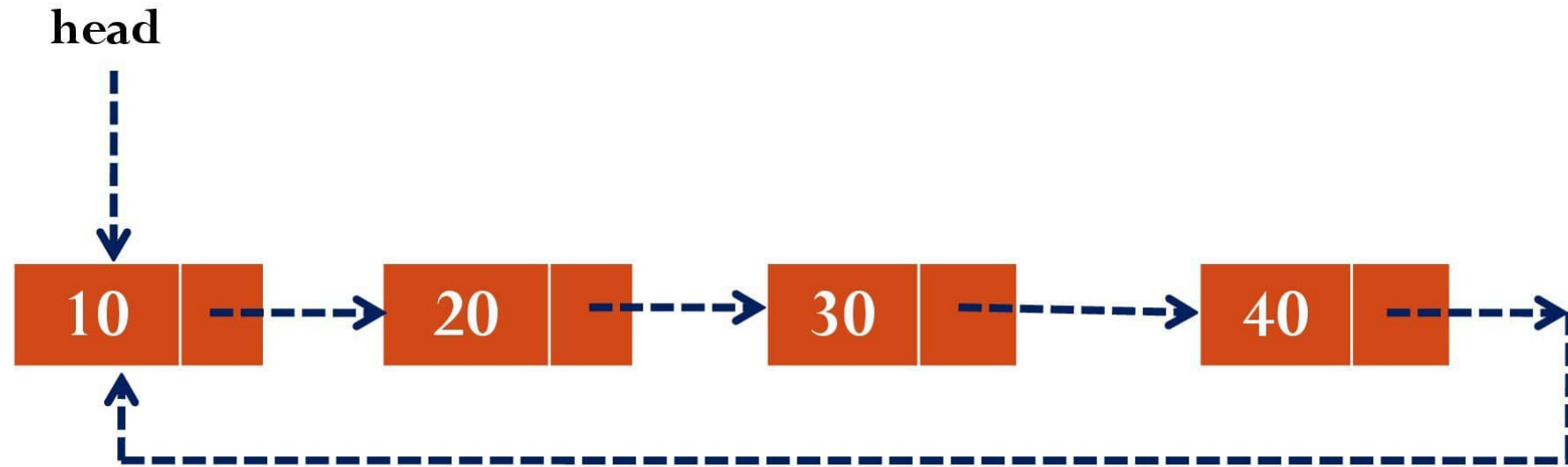
Case 2

Delete from End



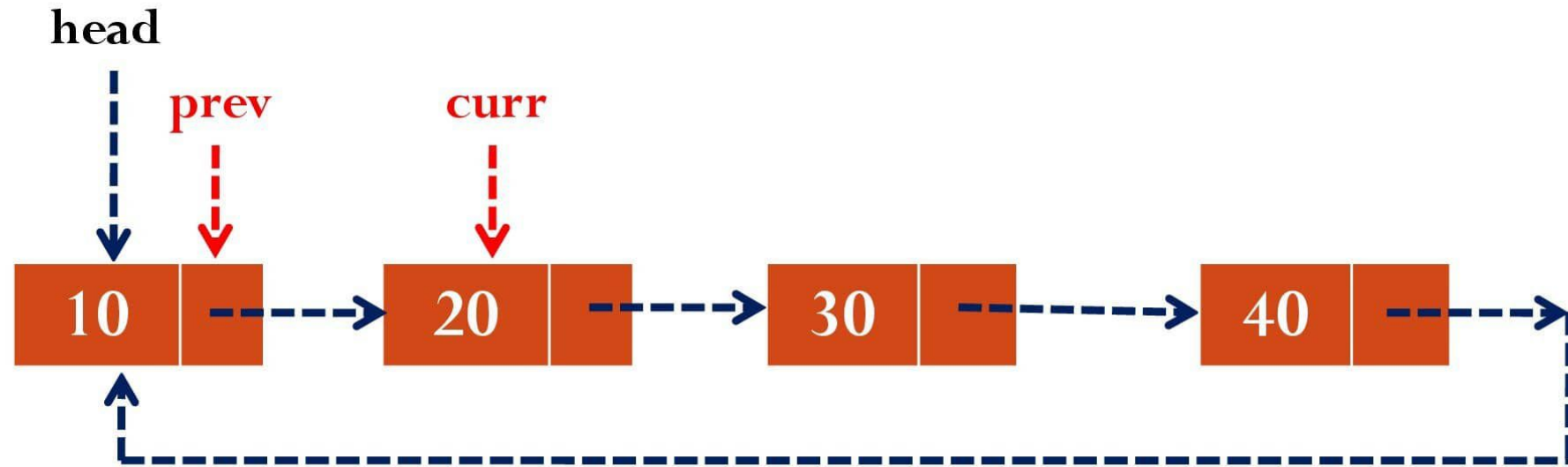
Case 3

Delete from End



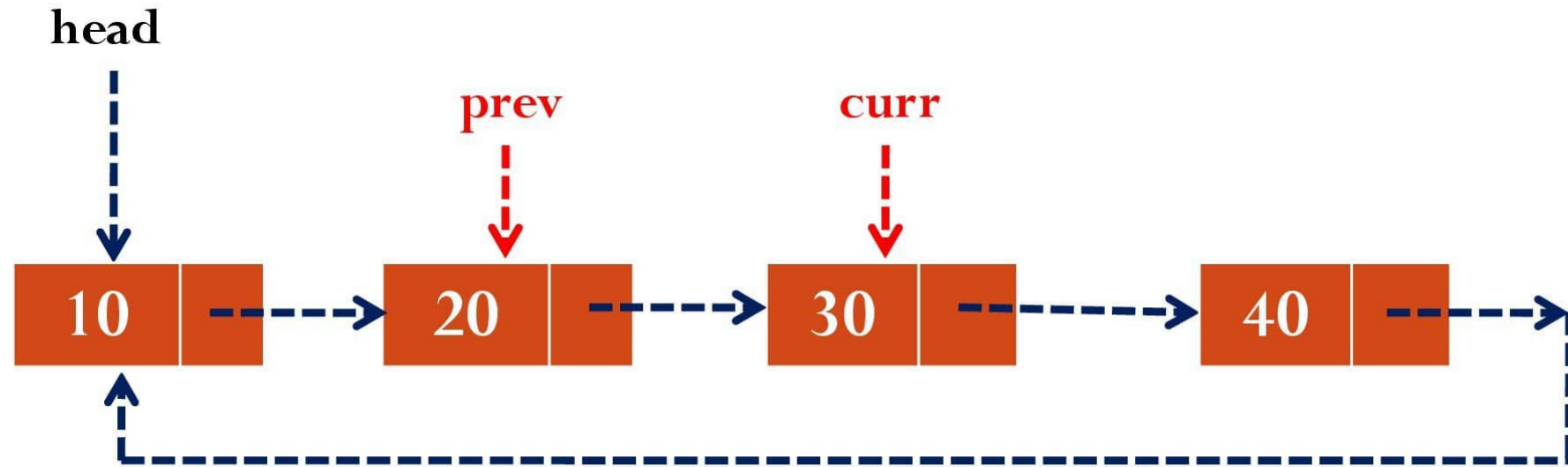
Case 3

Delete from End



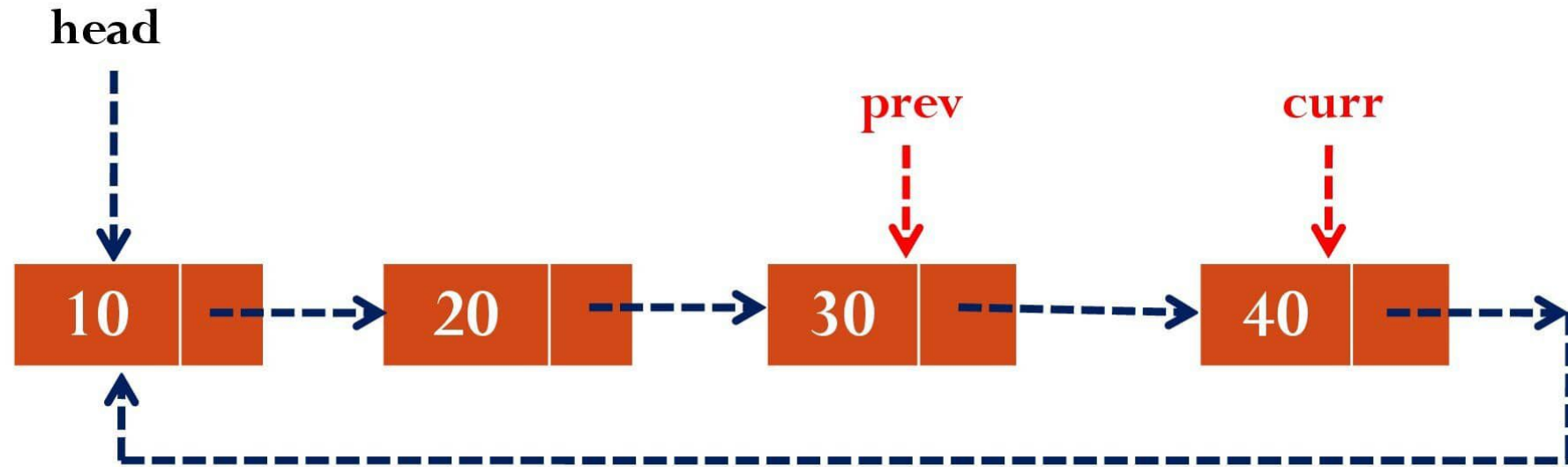
Case 3

Delete from End



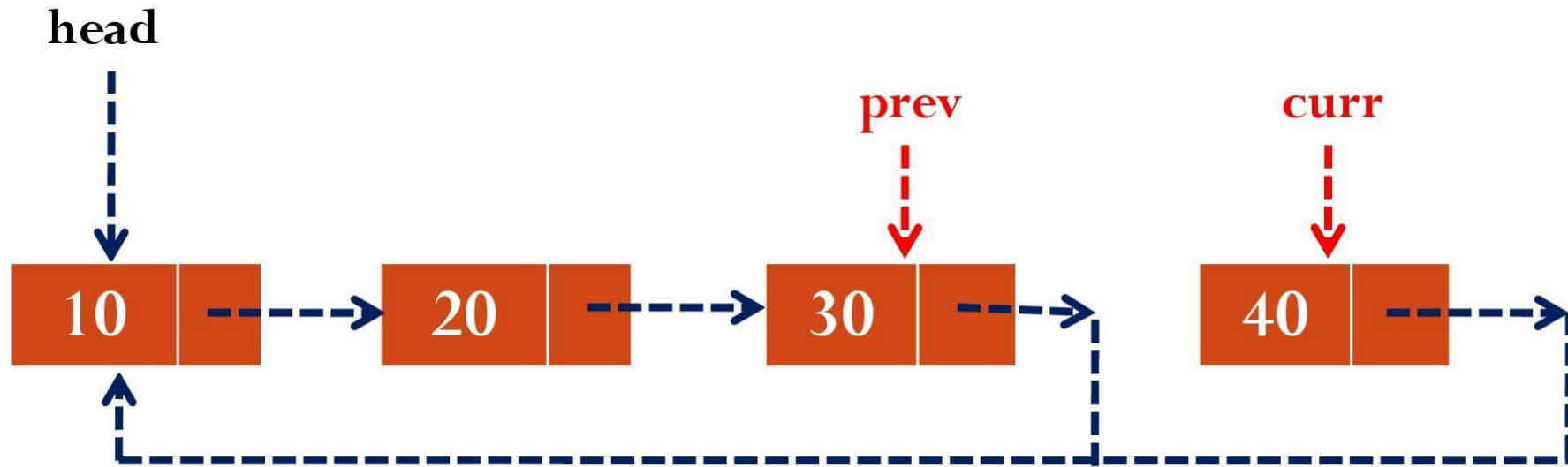
Case 3

Delete from End



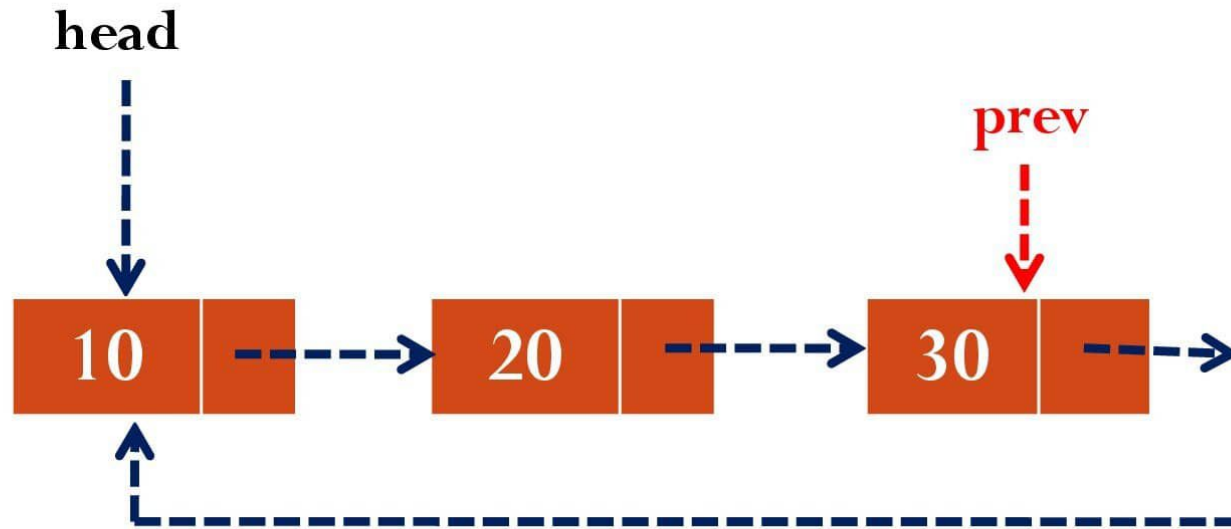
Case 3

Delete from End



Case 3

Delete from End



Delete from End- Algorithm

Algorithm Delete_End(head)

1. If head=NULL then
 1. Print “List is Empty. Deletion is not possible”
2. Else if head→link=head then
 1. temp=head
 2. head=NULL
 3. dispose(temp)
3. Else
 1. prev=head
 2. curr=head→link
 3. While curr→link!=head do
 1. prev=curr
 2. curr=curr→link
 4. prev→link=head
 5. dispose(curr)

Deletion

1. Delete from Front
2. Delete from End
3. Delete a specified node

Delete a specified node

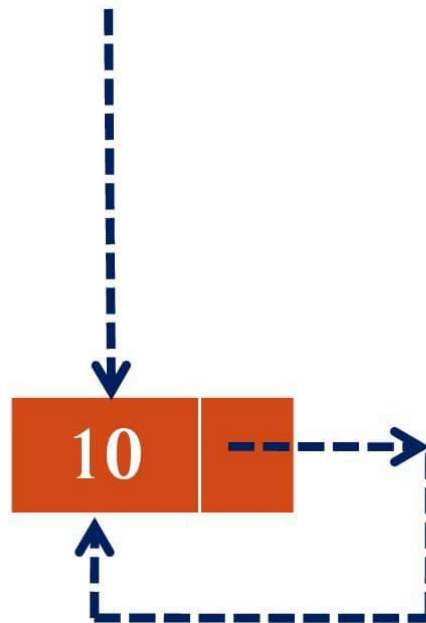
2 cases:

1. List is empty
2. List contains only one node
3. Search data is in the first node of the list
4. All Other cases

Case 2

Delete 10

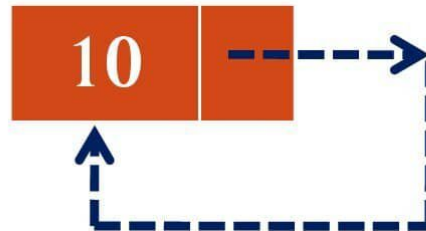
head



Case 2

Delete 10

head



Case 2

Delete 10

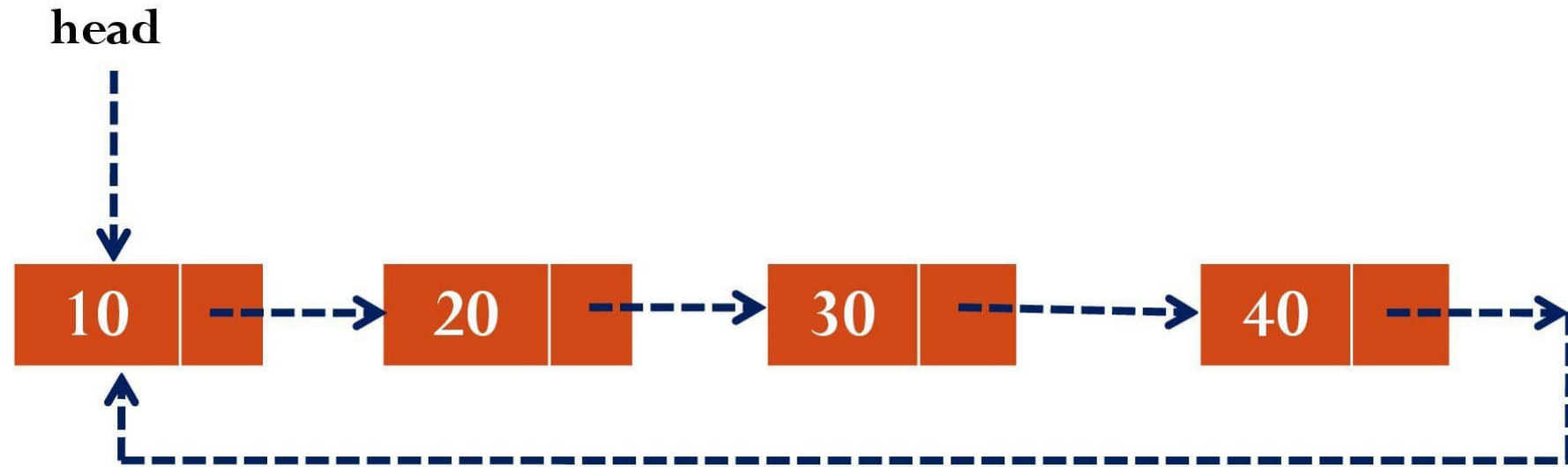
head



NULL

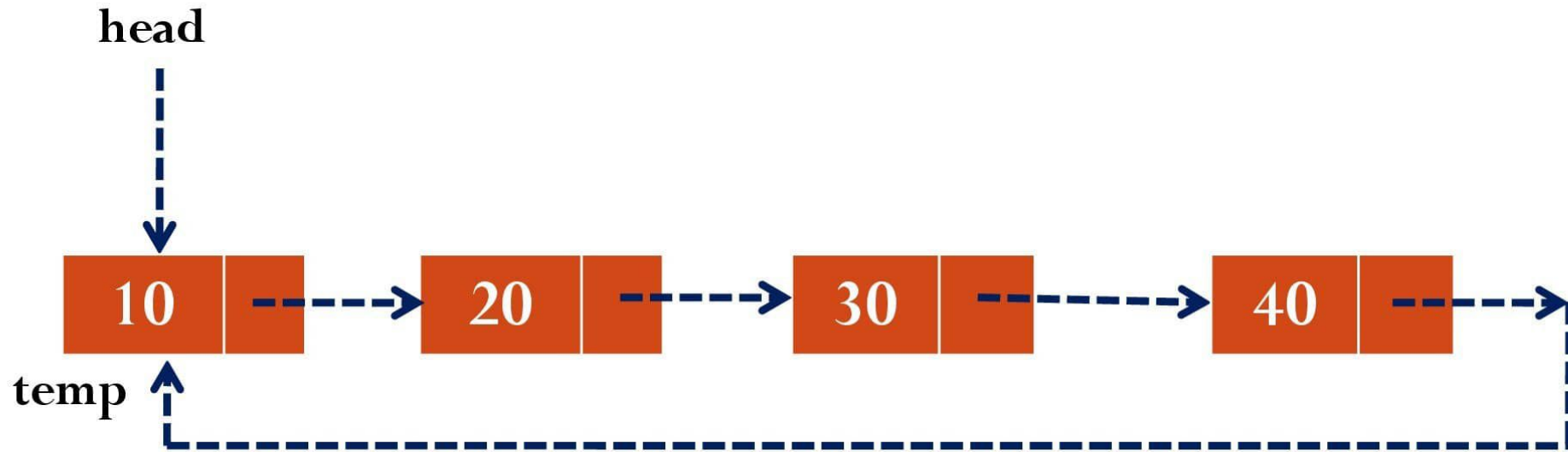
Case 3

Delete 10



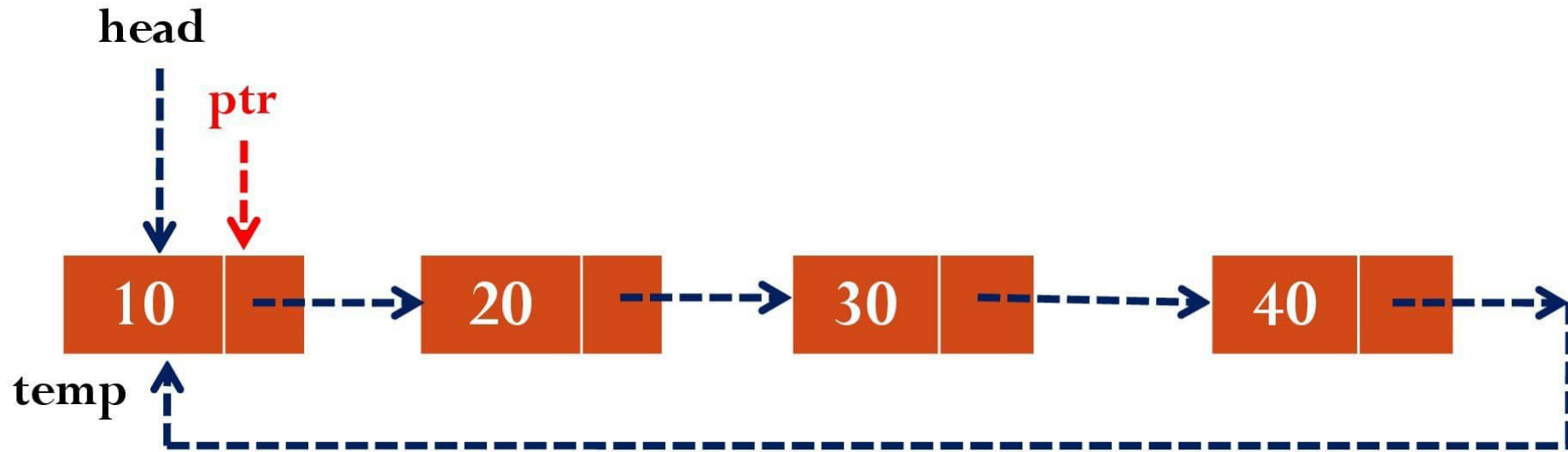
Case 3

Delete 10



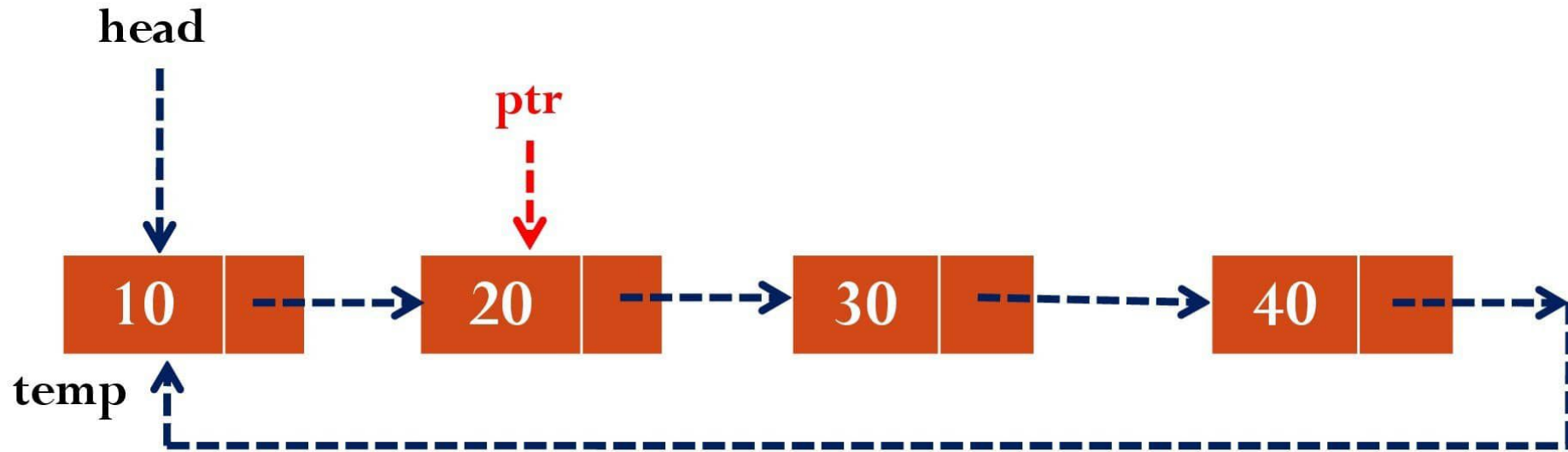
Case 3

Delete 10



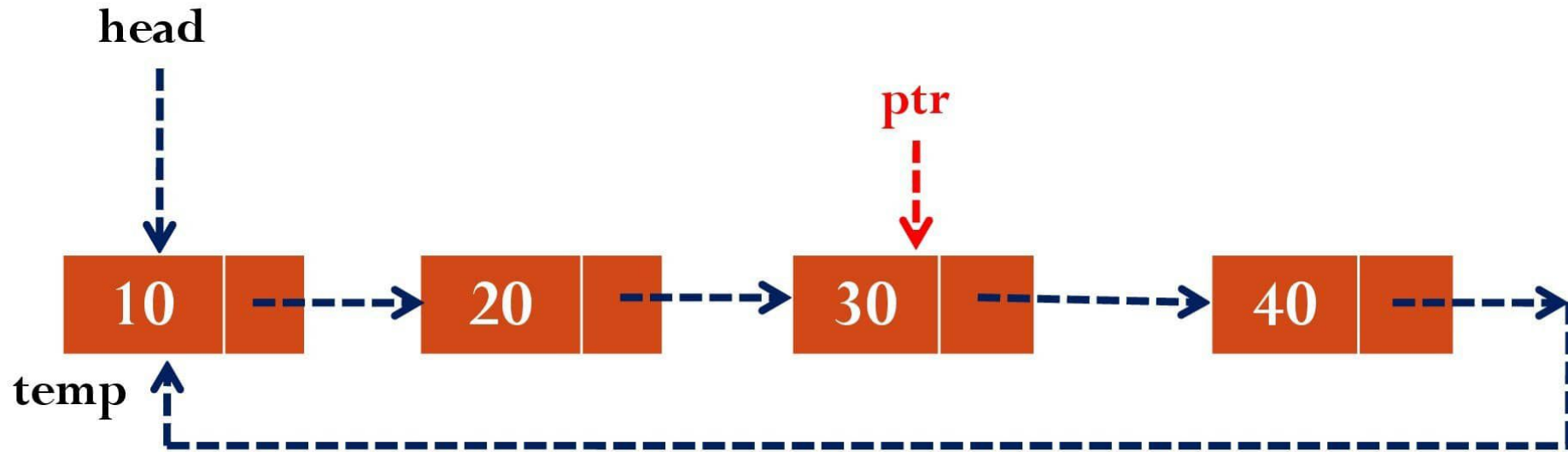
Case 3

Delete 10



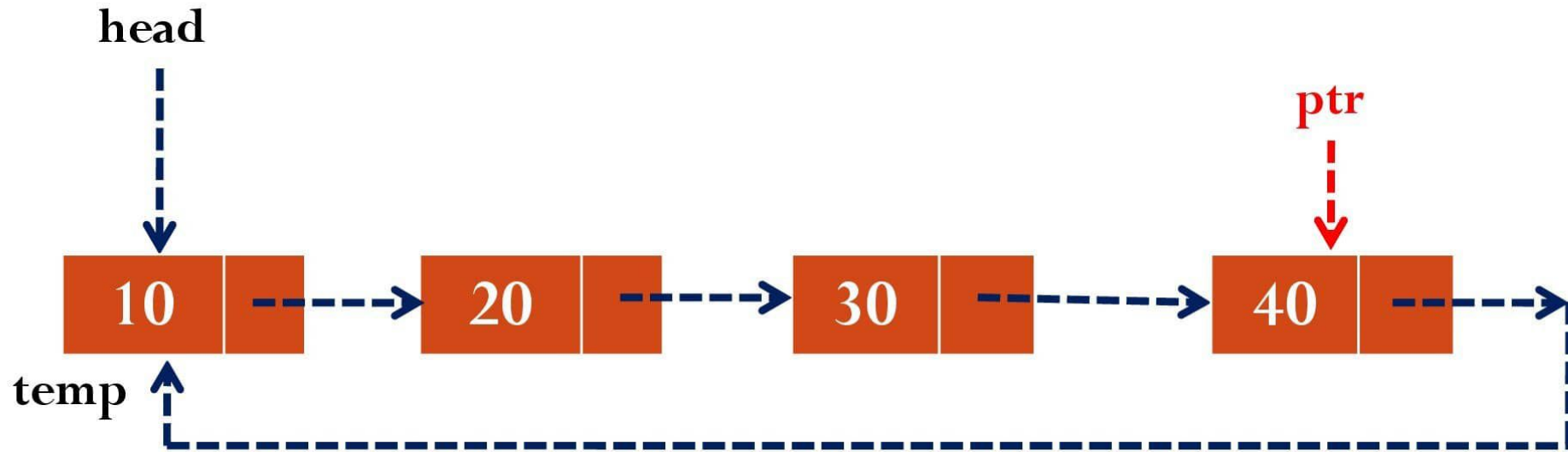
Case 3

Delete 10



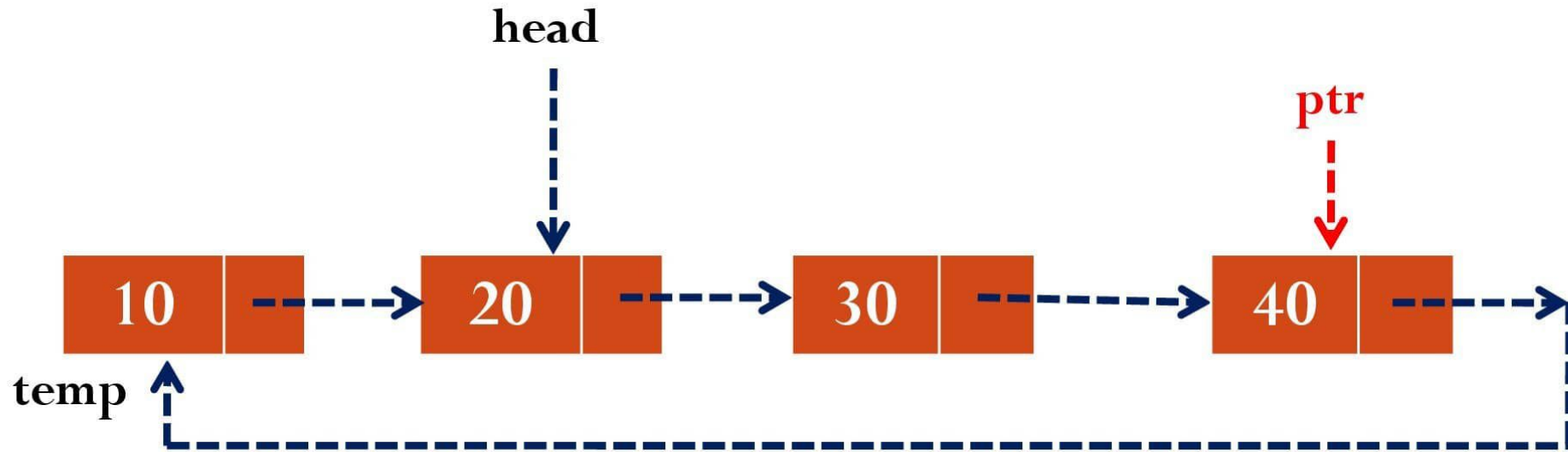
Case 3

Delete 10



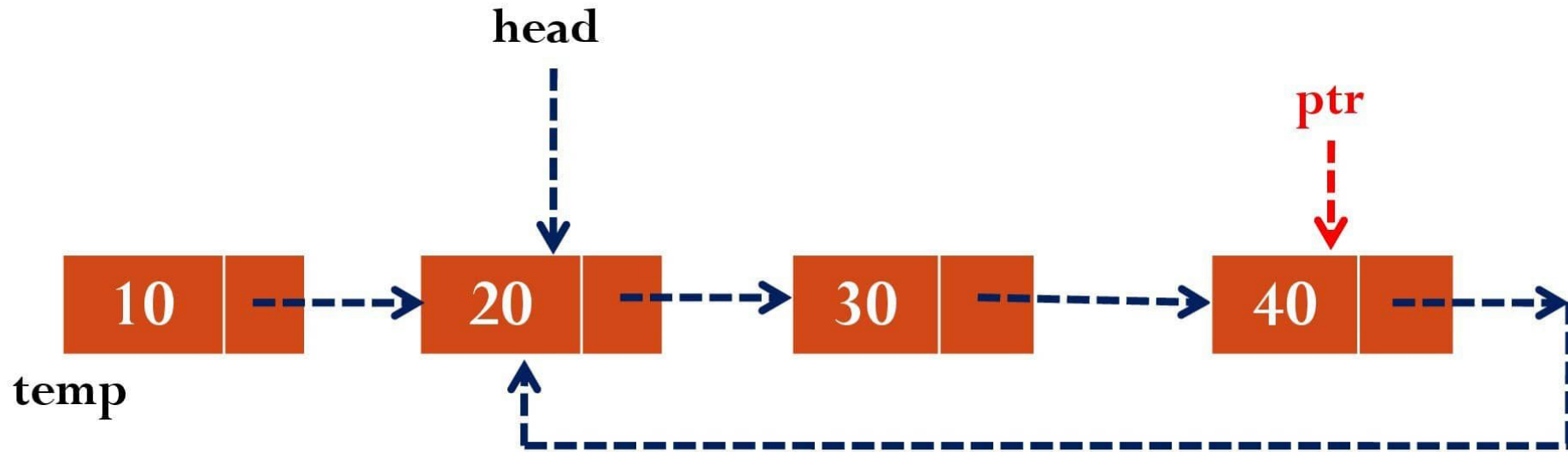
Case 3

Delete 10



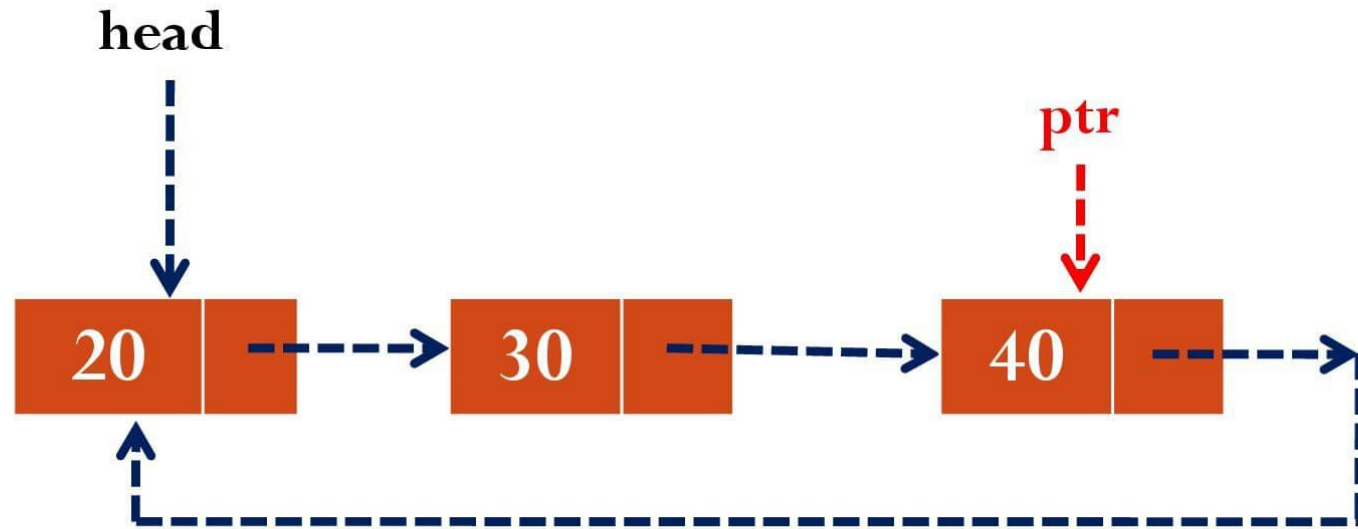
Case 3

Delete 10



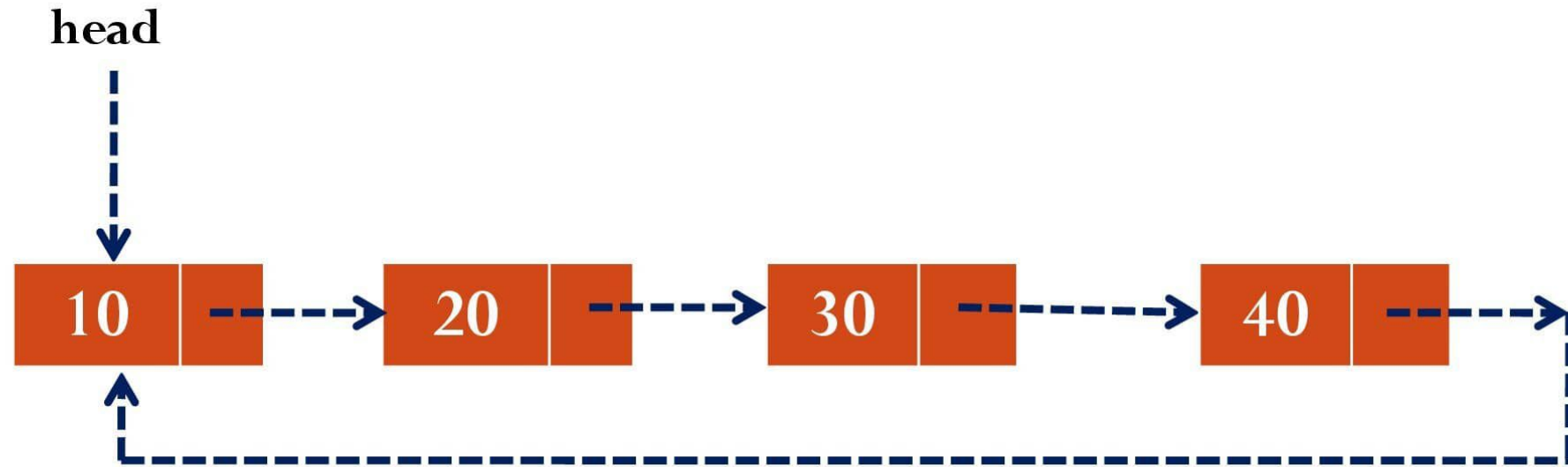
Case 3

Delete 10



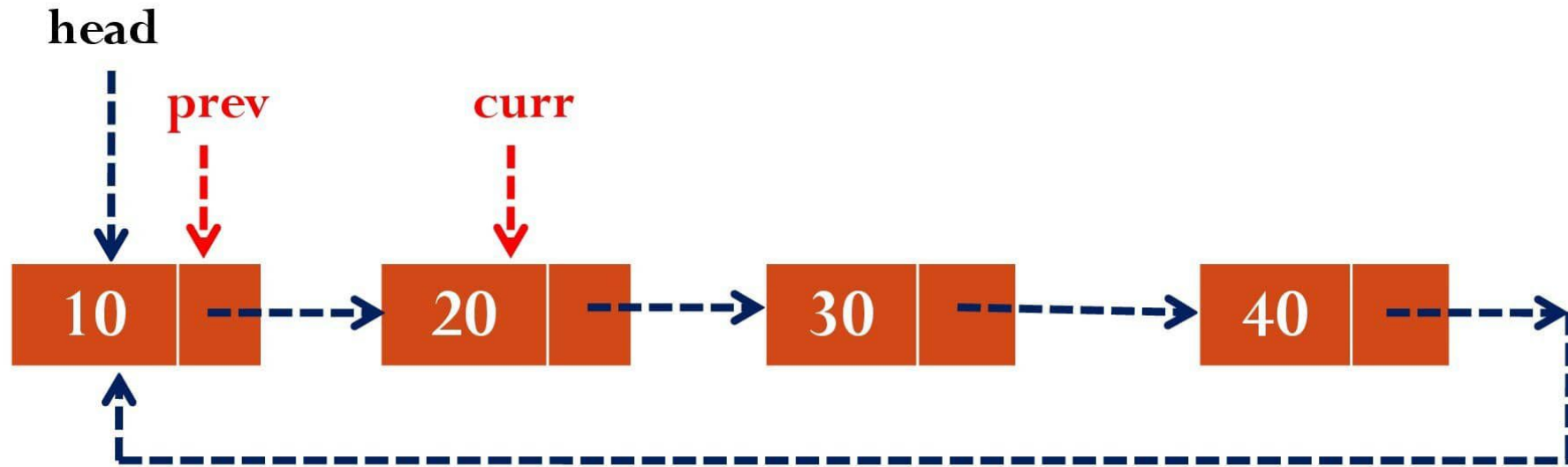
Case 4

Delete 30



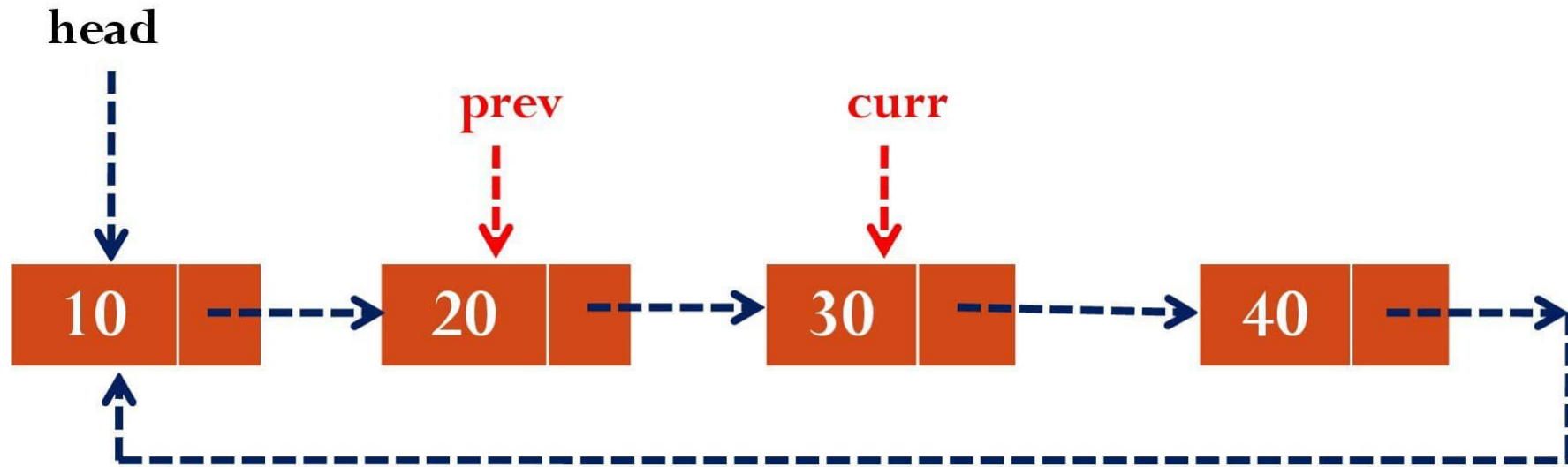
Case 4

Delete 30



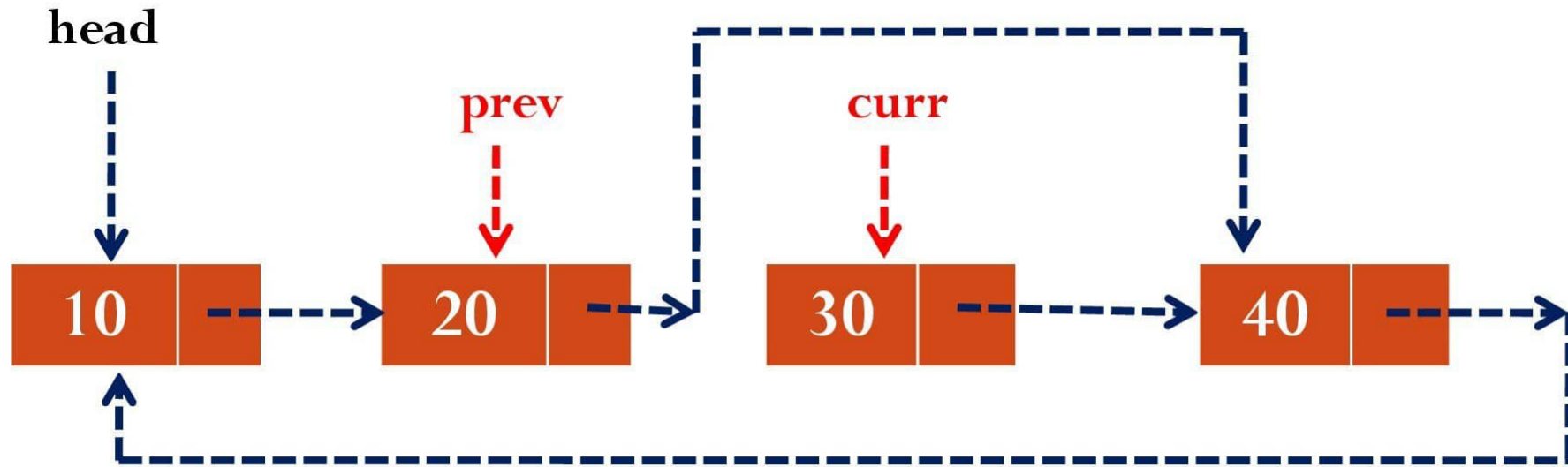
Case 4

Delete 30



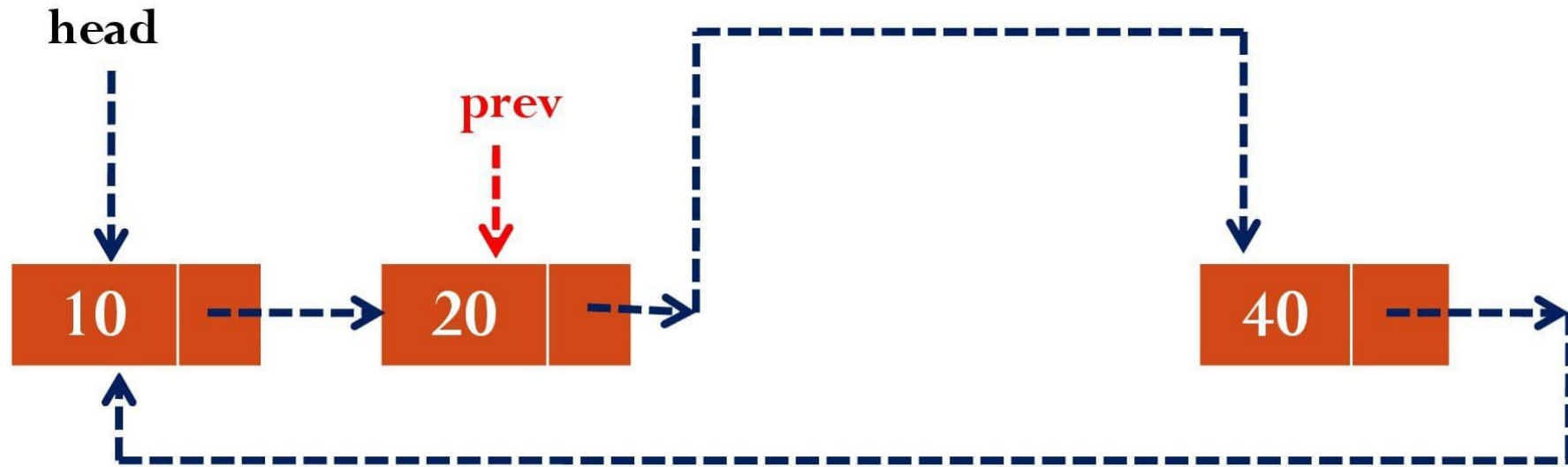
Case 4

Delete 30



Case 4

Delete 30



Delete a specified node~ Algorithm

Algorithm Delete_Any(head,key)

1. If head=NULL then
 1. Print “List is Empty. Deletion is not possible”
2. Else If head→link=head then
 1. If head→data=key then head=NULL
 2. Else Print “Search data not found.”
3. Else if head→data=key then
 1. temp=ptr=head
 2. While ptr→link!=head do
 1. ptr=ptr→link
 3. head=head→link
 4. ptr→link=head
 5. dispose(temp)

Delete a specified node~ Algorithm

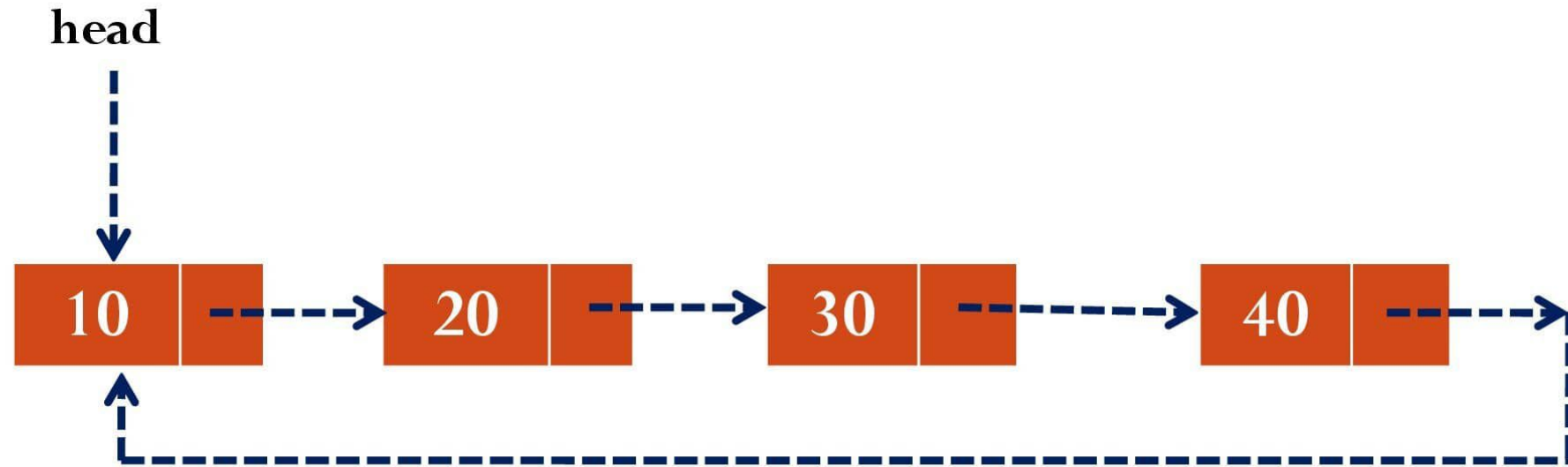
4. Else
 1. prev=head
 2. curr=head→link
 3. While curr→link!=head and curr→data!=key do
 1. prev=curr
 2. curr=curr→link
 4. if curr→data=key then
 1. prev→link=curr→link
 2. dispose(curr)
5. Else
 1. Print “Search data not found”

Search

1. List is empty
2. List is not empty

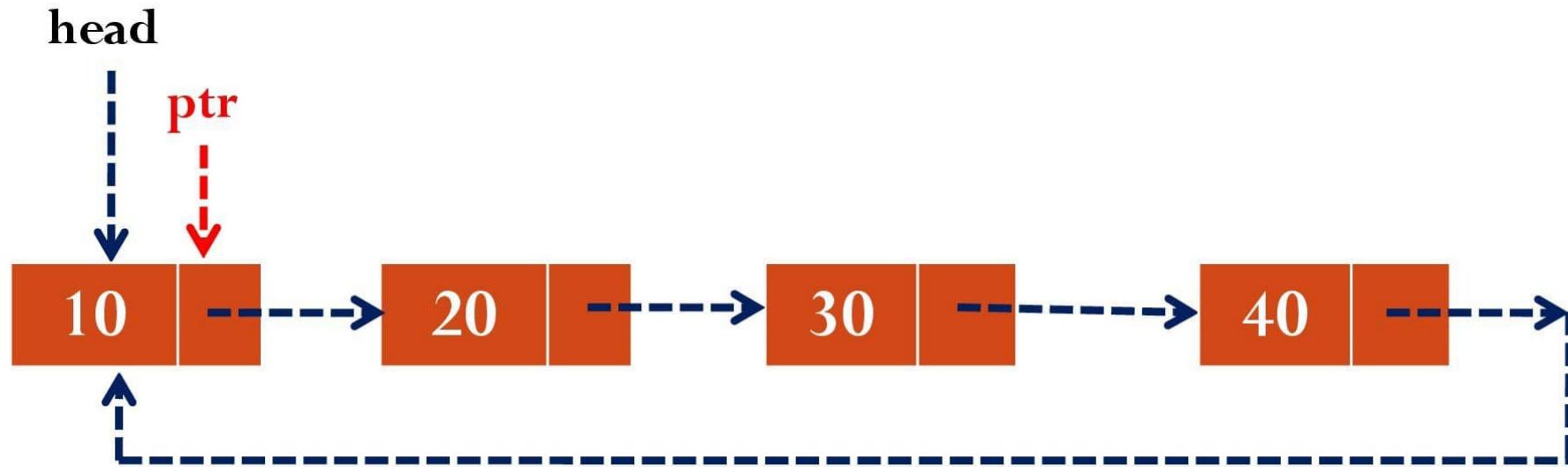
Case 2

Search 30



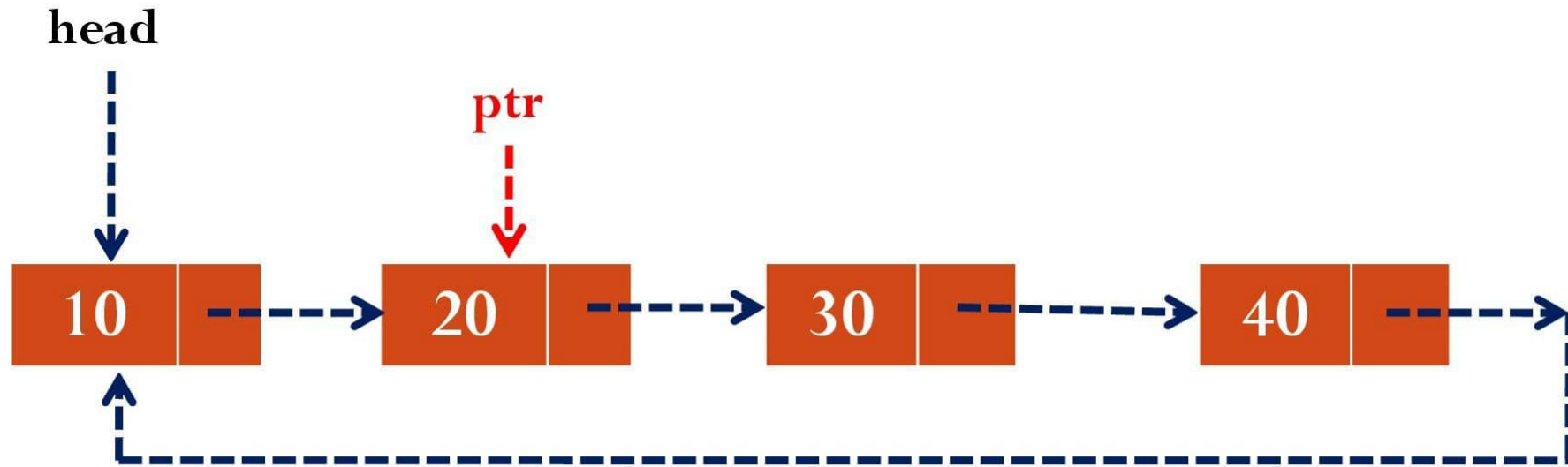
Case 2

Search 30



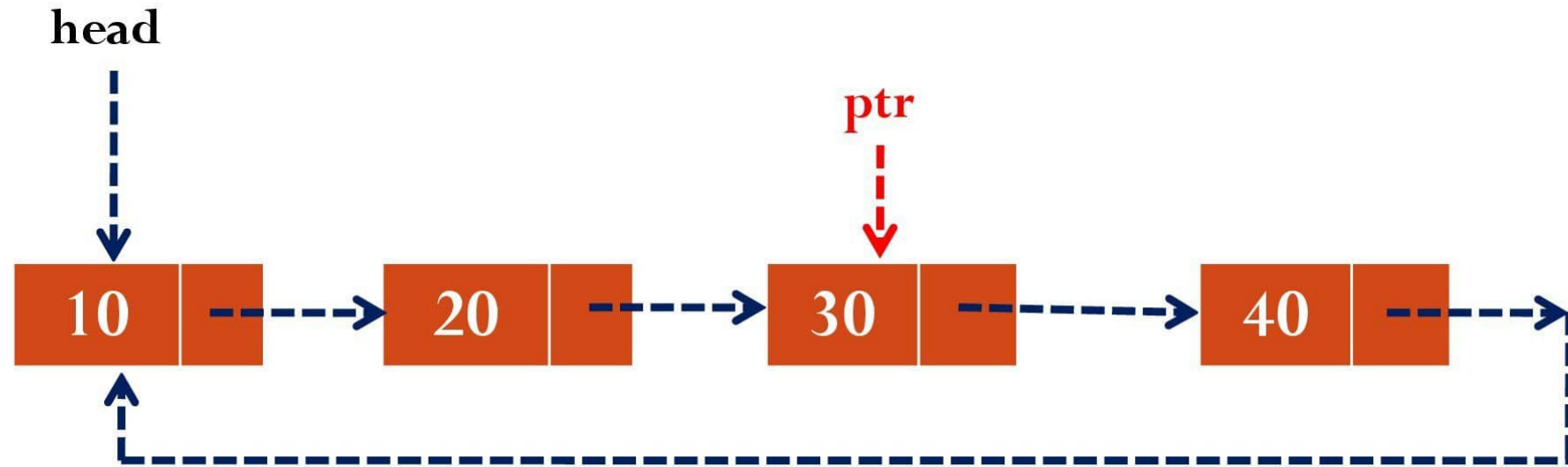
Case 2

Search 30



Case 2

Search 30

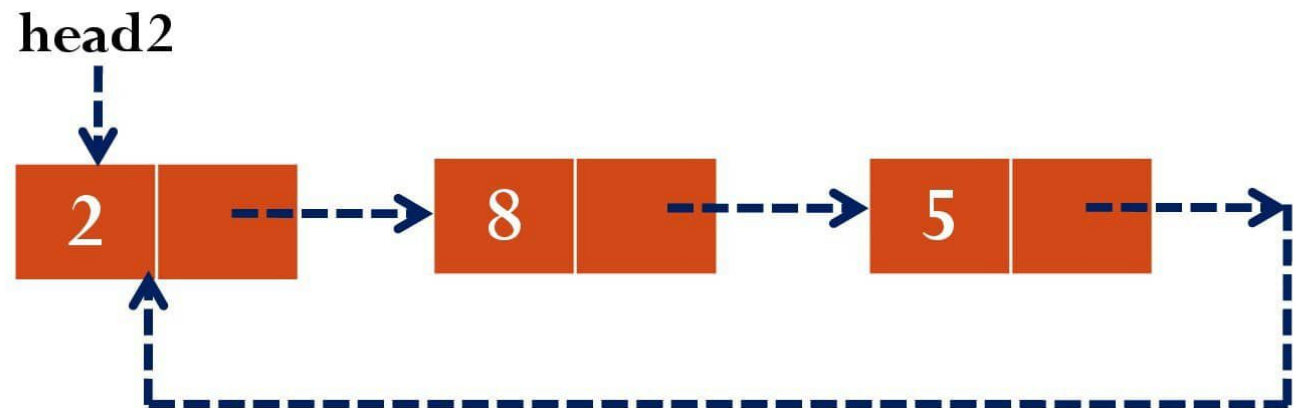
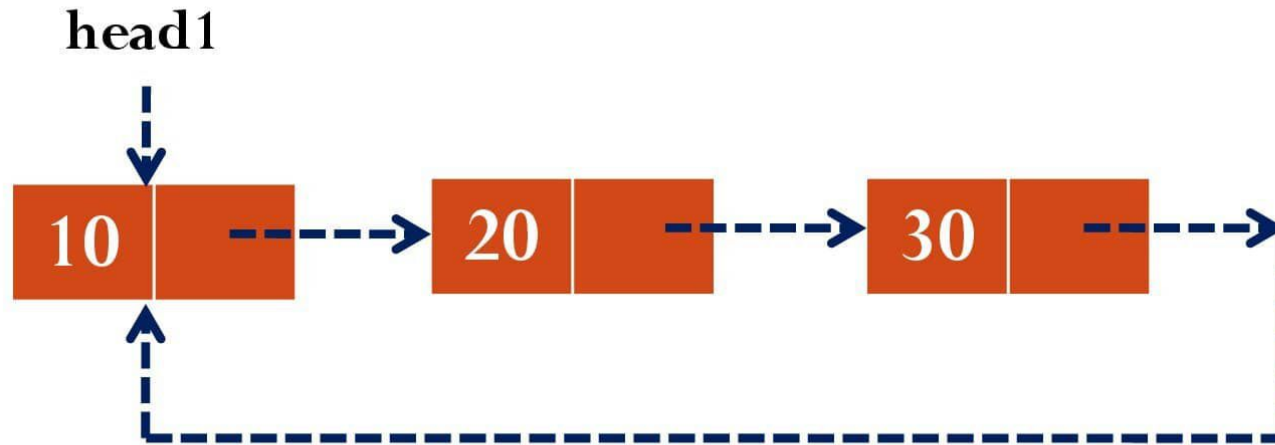


Search ~ Algorithm

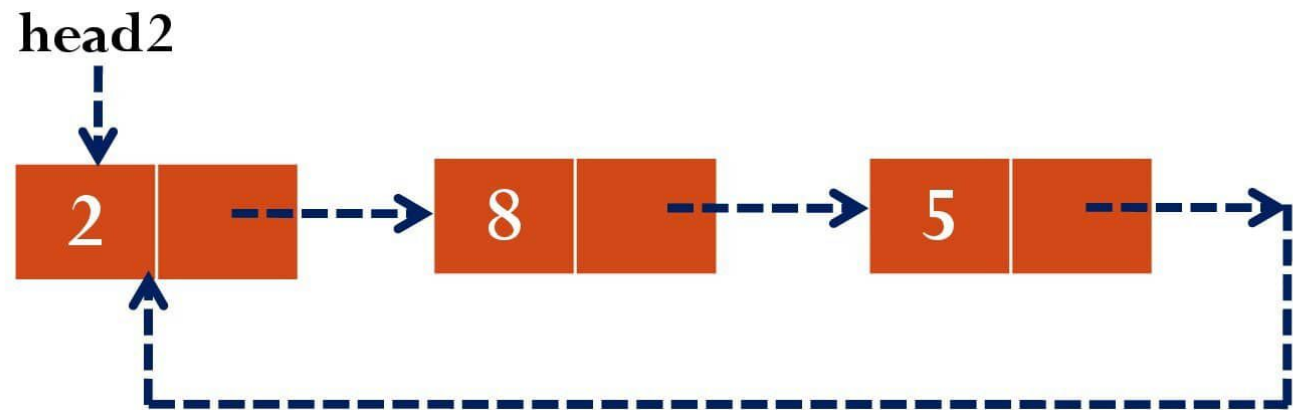
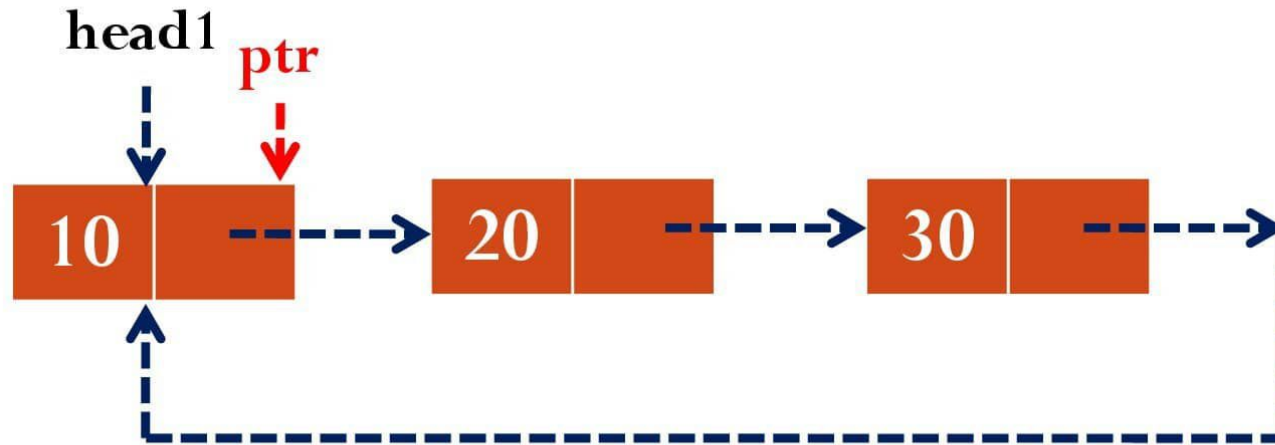
Algorithm Search(head, key)

1. If head=NULL then
 1. Print “List is Empty”
2. Else
 1. ptr=head
 2. While ptr→data!=key do
 1. ptr=ptr→link
 2. If ptr==head then
 1. break
 3. If ptr→data==key then
 1. Print “Search data found”
 4. Else Print “Search data not found”

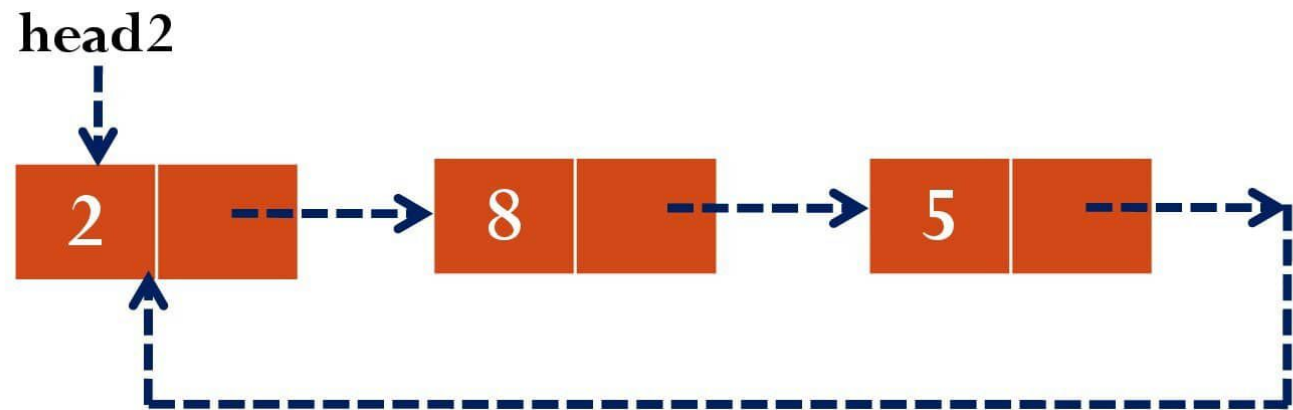
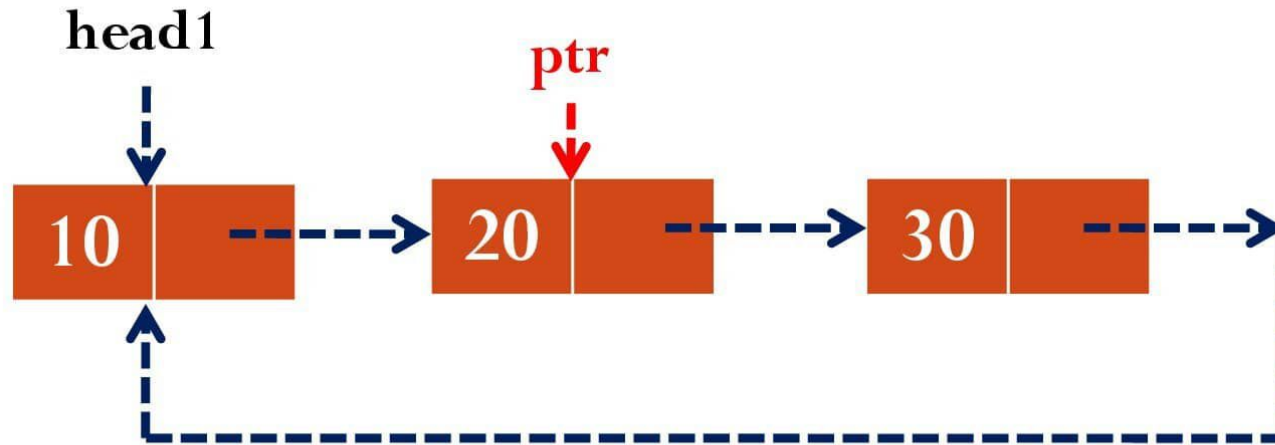
Merge



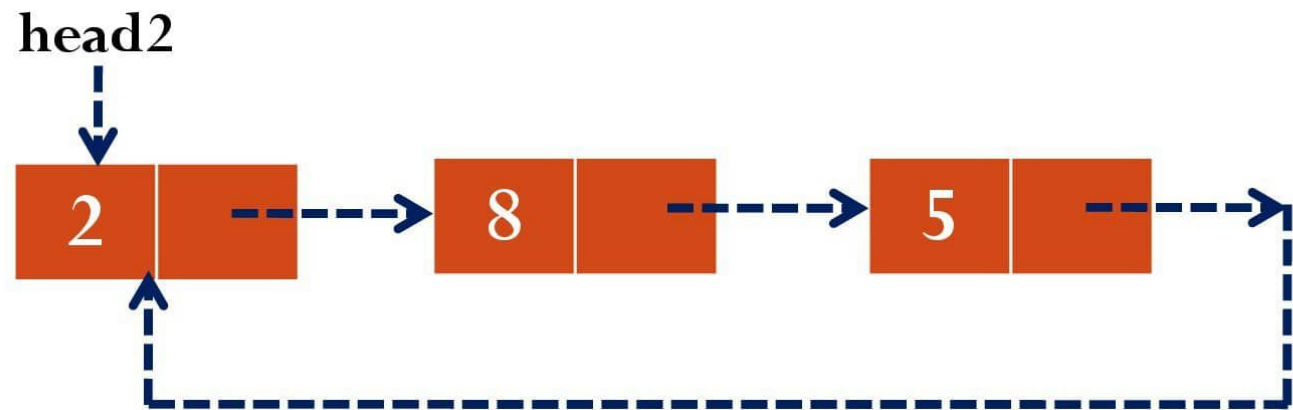
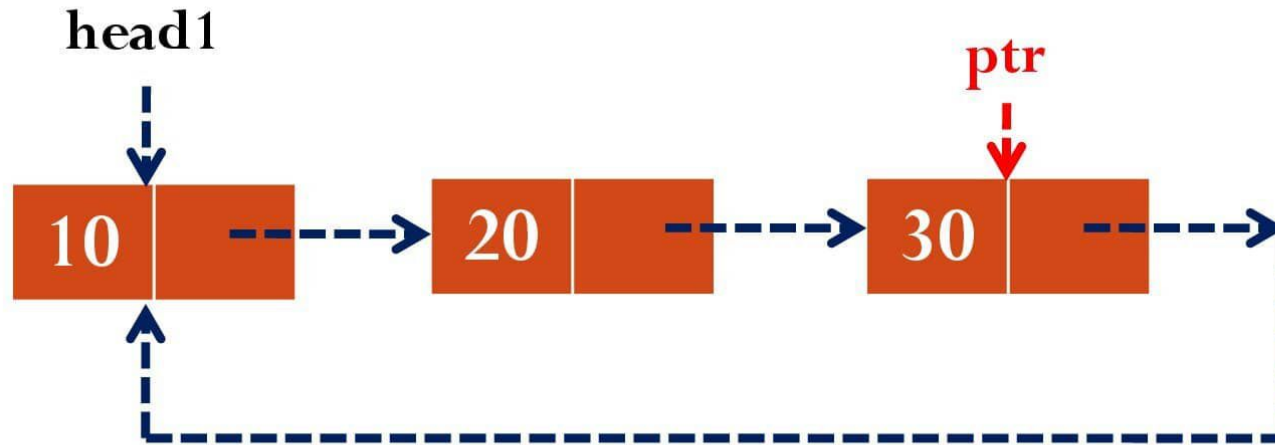
Merge



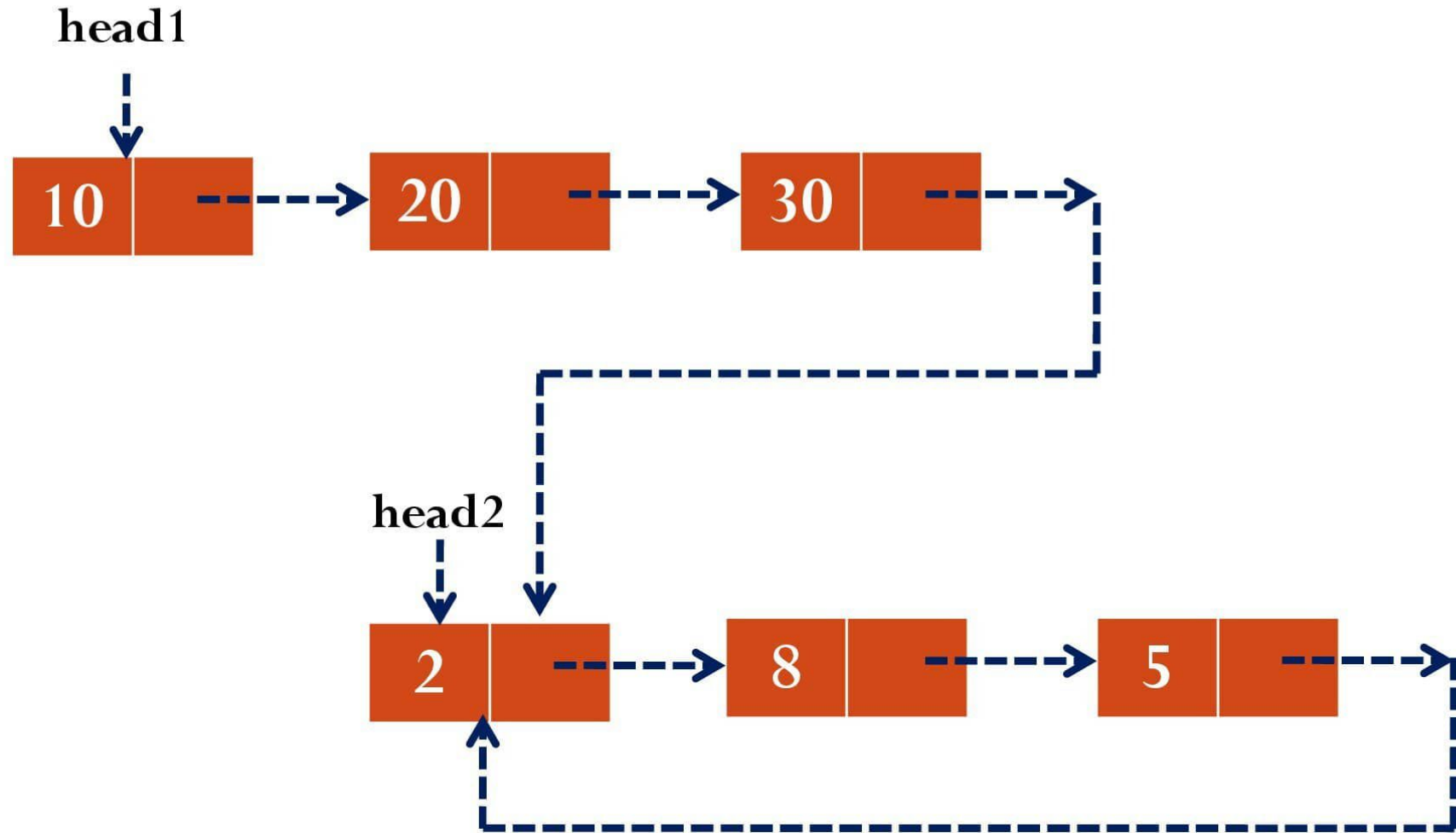
Merge



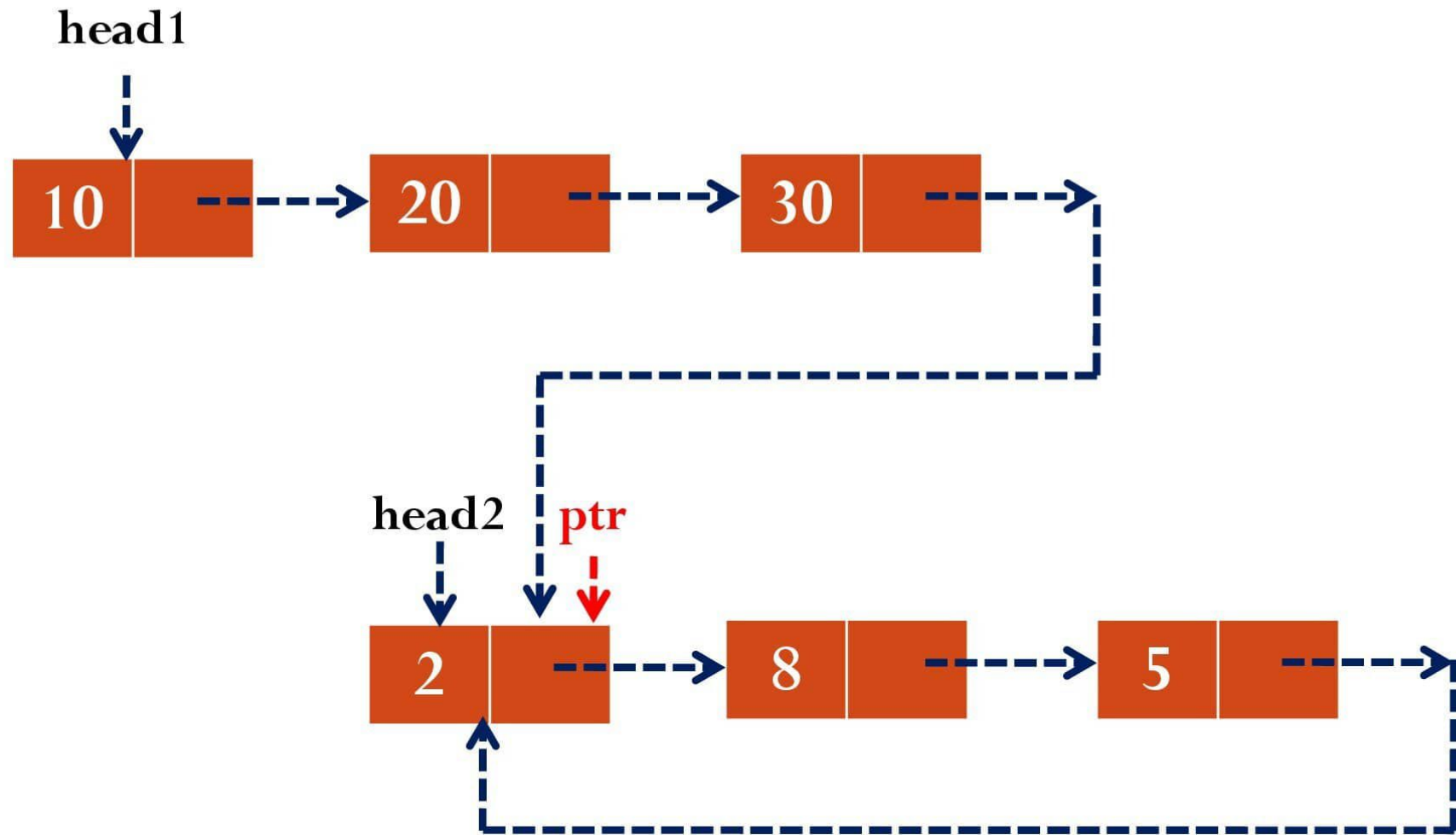
Merge



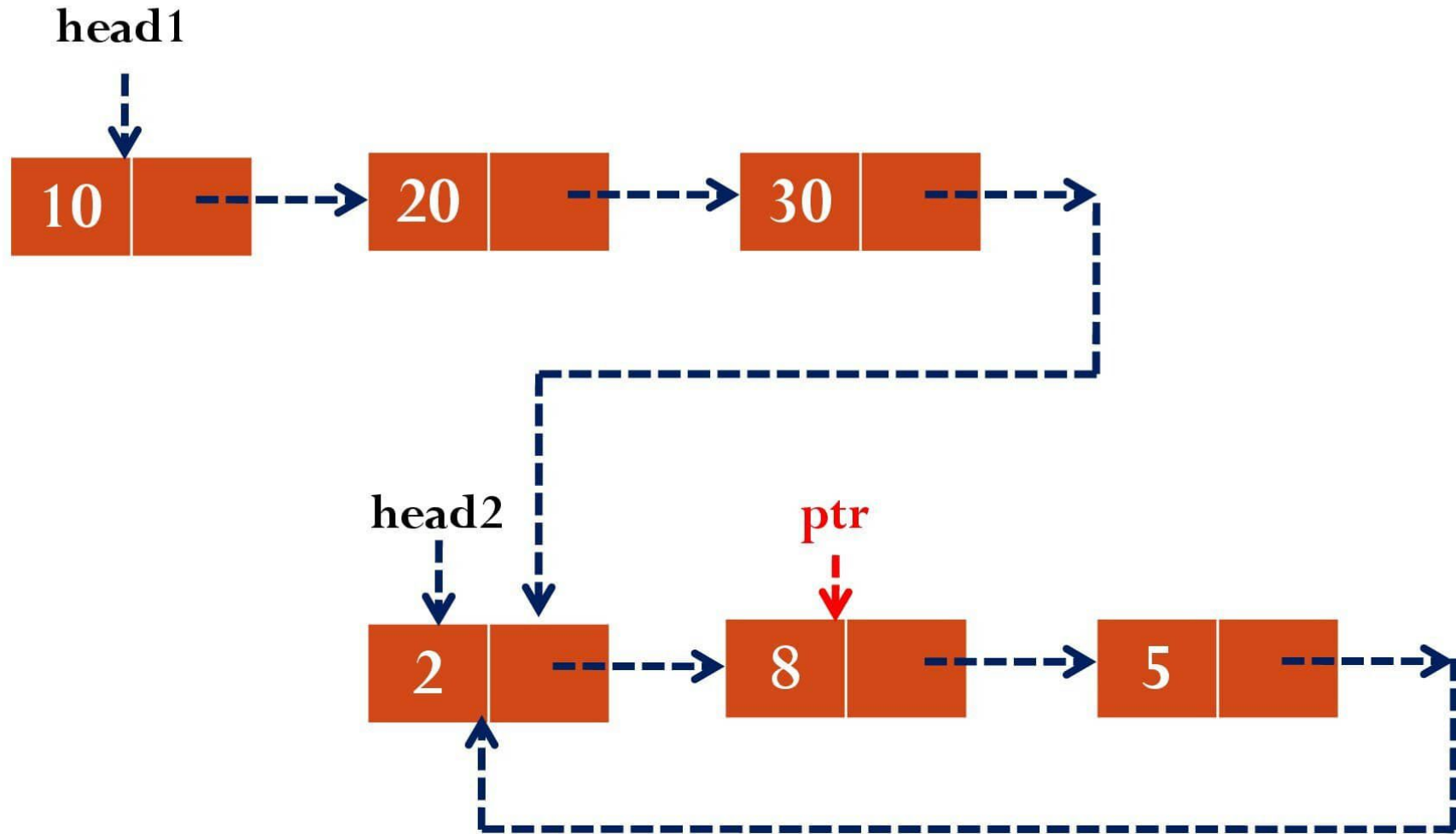
Merge



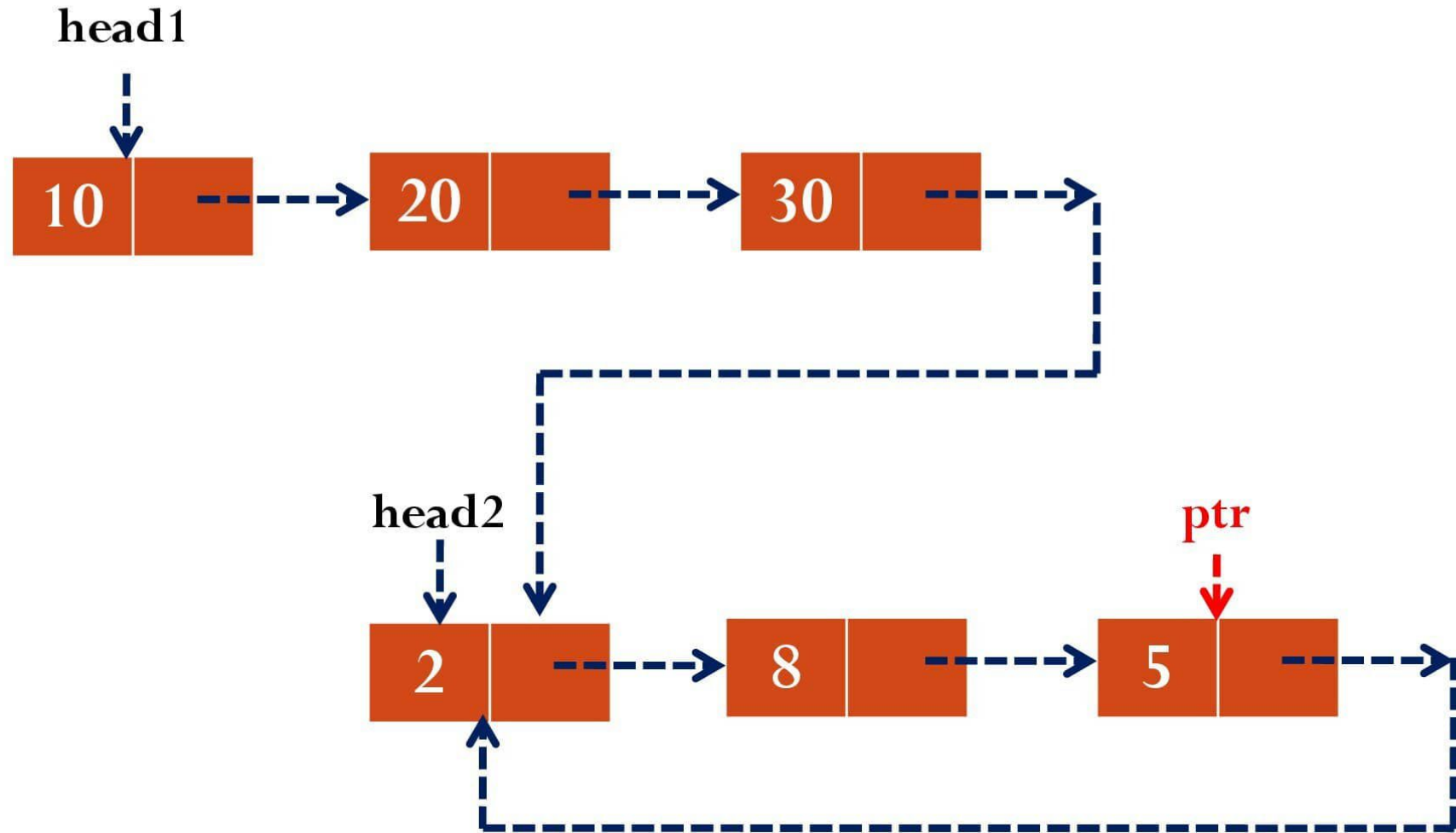
Merge



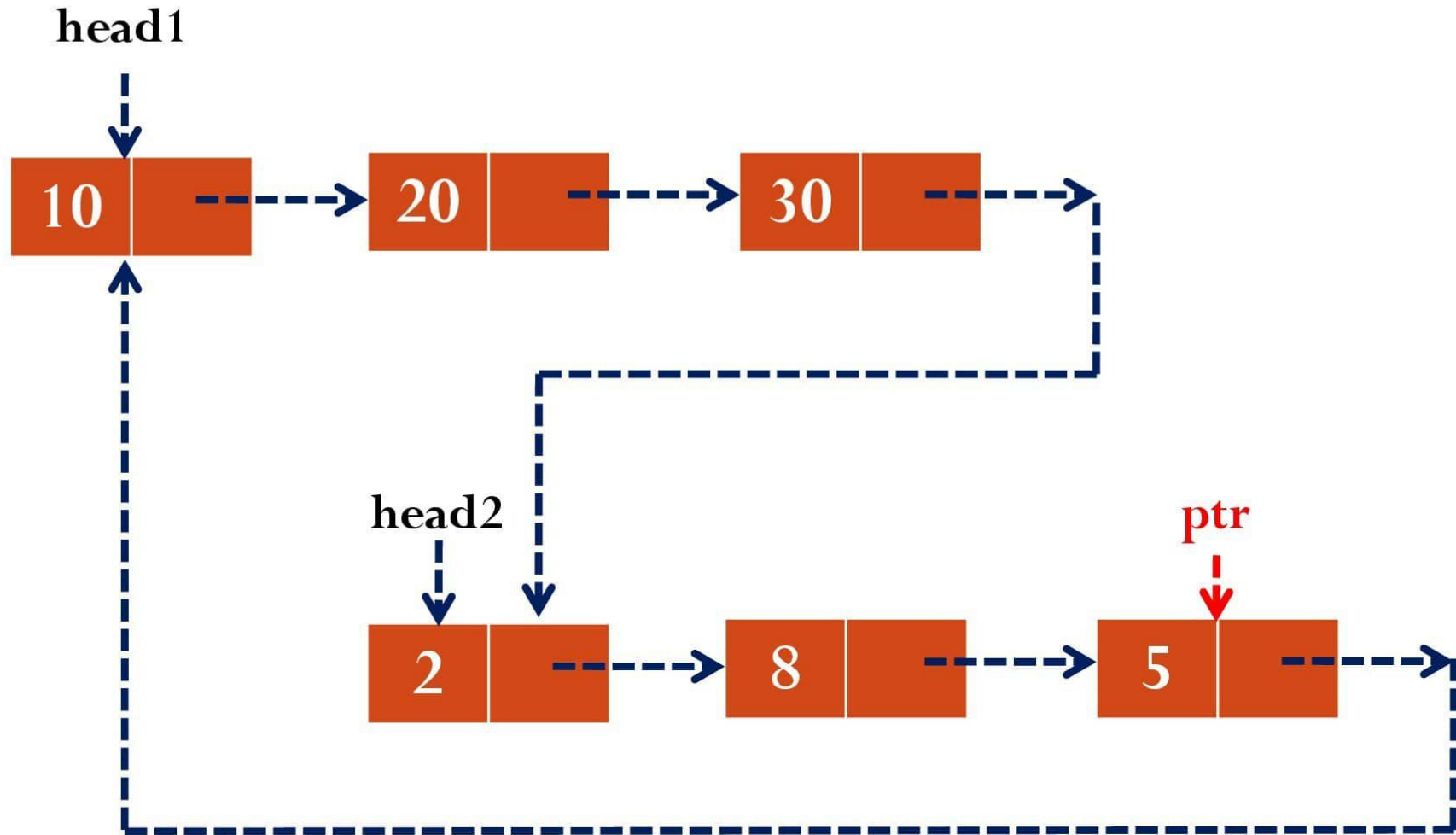
Merge



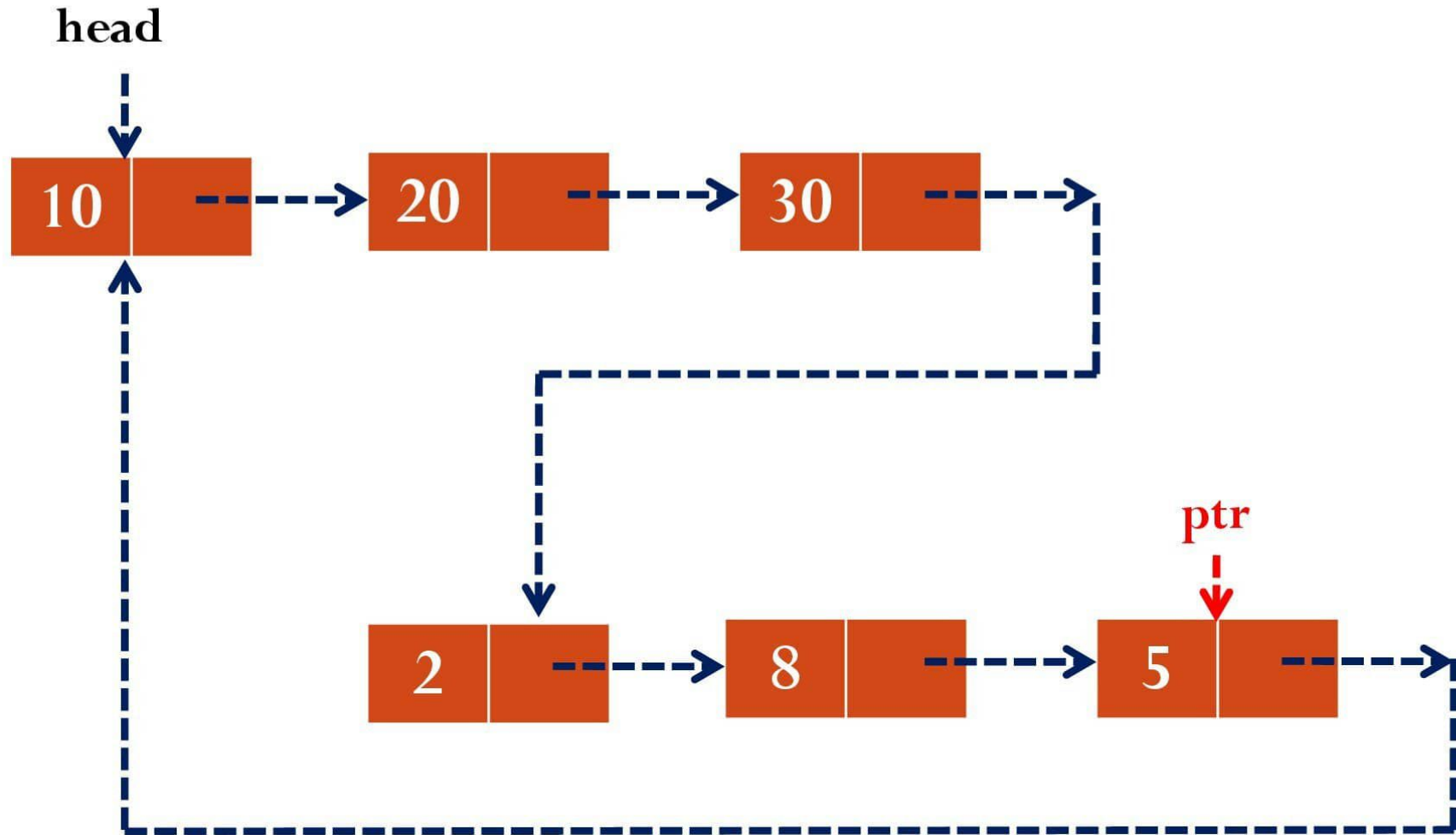
Merge



Merge



Merge



Merge ~ Algorithm

Algorithm Merge(head1, head2)

1. ptr = head1
2. while ptr → link != head1 then
 1. ptr = ptr → link
3. ptr → link = head2
4. ptr = head2
5. while ptr → link != head2 then
 1. ptr = ptr → link
6. ptr → link = head1
7. head = head1